# WD16 MICROCOMPUTER

(Using MCP 3-Chip Microprocessor Set)

**PROGRAMMER'S REFERENCE MANUAL** 



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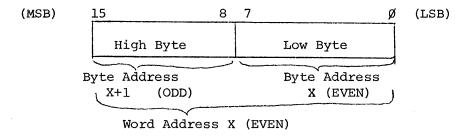
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#### CHAPTER 1 - GENERAL

The WD1600 microcomputer is a 16 bit machine with both word and byte addressing, an automatic push down hardware stack, vectored interrupt handling, eight 16 bit registers, and PC relative addressing. A byte is defined as 8 bits, and a word is defined as 2 bytes. A memory address increment of one is an increment of 1 byte. An address increment of two is an increment of 1 word. Word addresses always start on even bytes. For any memory location the even byte is the least significant byte. Bit 0 is defined as the LSB of a memory location.



Unless otherwise stated, word addressing is implied. All addresses and op codes are done in hex unless otherwise stated. All hex numbers are enclosed within double quotes.

## LEGEND OF ABBREVIATIONS

REG = Register

SRC = Source Address

(SRC) = Contents of Source Address

DST = Destination Address

(DST) = Contents of Destination Address

 $(SRC)_B$  = Contents of Source Byte Address

 $(DST)_B$  = Contents of Destination Byte Address

 $\overline{x}$  = Ones Complement of X

 $\rightarrow$  = Twos Complement of X

 $\Delta$  = Logical And

∇ = Logical Or

@ = Indirect

↓ = Push

 $\uparrow$  = Pop

← = Destination Direction

+ = Addition

- = Subtraction

\* = Multiplication

/ = Division

: = Double Precision Chain Link

## PROCESSOR STATUS WORD

A 16 bit Processor	Status (PS)	Word	exists.	The format	is as follows:
	7				
Ext. Status Reg.	ALU	N	ZVC		

Where bits 8-15 are the contents of the external status register (see chapter 2), bits 4-7 are the status of the microprocessor ALU flags, and bits  $\emptyset$  -3 are the status of the condition indicators at the time the PS is formed. The ALU flags are of no use or concern to the programmer. They are stored along with the condition indicators automatically as a function of the micro-op. The four condition flags are updated during the execution of most op codes, and are used by the branch instructions to test for valid branch conditions. The exact status of each indicator is defined along with the descriptions of individual op codes in chapter 3. In general, however, the indicators are set by the following conditions:

N = set if the MSB of the result is set.

Z = set if the result is zero.

V = set if arithmetic overflow (underflow) occurs during addition (subtraction). Set to exclusive -or of N and C indicators otherwise.

C= set if carry (borrow) occurs during addition (subtraction). Also set to last bit shifted out during a shift operation.

## REGISTERS

There are 8 registers in the WD1600. All are 16 bits long. Six can be used as either accumulators or index registers, one is the stack pointer (SP), and one is the program counter (PC). The registers are numbered RØ - R7 with R6 = SP and R7 = PC. The register set is usually referred to in the following manner: RØ - R5, SP, PC.

## CHAPTER TWO - INTRODUCTION

#### ADDRESSING MODES

In general there are 8 addressing modes for both source and destination addressing. Not all op codes accept all 8 modes (see chapter 3). Those that do use the following format: 3 bits for the index register (RØ - R5, SP, PC) and 3 bits for the mode. The mode bits are the upper 3 bits of the 6 bit set. The modes are defined below. The numbers in parenthesis refer to notes that follow the definitions.

\_\_\_\_\_

MODE	NAME SY	MBOLIC	DESCRIPTION
ø	Direct Register	REG	REG is or contains operand.
1	Indirect Register	@REG	REG contains address of operand.
2	Auto-increment	(REG)+	REG contains address of operand. REG is post-incremented (1).
3	Auto-increment deferred	@(REG)+	REG contains address of add- ress of operand. REG is post- incremented by 2.
4	Auto-decrement	- (REG)	REG is predecremented (1). REG then contains address of operand.
5	Auto-decrement deferred	@-(REG)	REG is predecremented by 2. REG then contains address of address of operand.
6	Indexed register	X(REG)	Contents of REG plus X is address of operand (2).
7	Indexed register deferred	@X(REG)	Contents of REG plus X is address of address of operand (2).

\_\_\_\_\_

NOTE 1: For word operations the increment/decrement is 2. For byte operations the increment/decrement is 1 unless the index register is SP or PC. In this case the increment/decrement is always 2.

NOTE 2: The contents of REG remain unchanged.

When using PC as the index register the assembler accepts the following 4 formats in place of the formats mentioned above for ease of programming.

MODE	NAME	SYMBOLIC	DESCRIPTION
2	Immediate	#N	Operand N follows op code.
3	Absolute	@#N	Address of operand is N and it
			follows the op code in memory.
6	Relative	A	PC relative offset to address A, which contains operand, follows
			op code.
7	Relative deferred	A9	PC relative offset to address A, which contains address of operand,
			follows the op code.

The 8 modes are referred to as Source Mode Ø to Source Mode 7 (SMØ -SM7) and Destination Mode Ø to Destination Mode 7 (DMØ -DM7). In Chapter 3 these modes are referred to in general terms during op code definitions as "SRC" and "DST".

#### STACK OPERATIONS

Although automatic stack operations are provided for, no specific area of memory is set aside for the stack. The user must assign an area of memory by loading the stack pointer with the top address of the designated stack area. Stack operations are pushdown pop-up operations with predecrements and post-increments of SP. Stack operations may also be executed explicitly by using SP as an index register with op codes that allow SMØ - SM7 and/or DMØ - DM7 addressing.

When pushing the PS the word is formed just prior to the push. When popping the PS the condition indicators and interrupt enable flag are set to the status of the appropriate bits in the popped PS. Other than that the popped PS goes nowhere. Unless otherwise stated popping the PS from the stack performs the above mentioned operations and only the above mentioned operations.

When pushing the PC onto the stack PC will be set to the address of the op code that follows the op code that caused the push. There are cases where some op code formats can alter this rule. They generally involve advanced programming techniques. A few are mentioned in appendix C. In particular, system errors that are caused by programming errors and not real time error conditions will push a PC that points to the op code that follows the op code that caused the error. The stored PC must be decremented by two to get the address of the offending op code.

#### INTERRUPT LINES

There are 4 interrupt lines available to the system. They are labeled  $I\emptyset$  - I3. These lines are assigned functions as follows:

IØ = Vectored interrupt line

Il = Nonvectored interrupt line

I2 = Enable/disable for IØ and Il.

I3 = Halt switch

The priority among the lines is as follows:

I3, I1 $\Delta$ I2, I $\emptyset$  $\Delta$ I2.

Note that I3 is always enabled. Note also that the nonvectored interrupt has priority over the vectored interrupt. The system is currently set up so that power fail and a real time clock can be assigned to I1, and up to 16 devices assigned to IØ.\* The two interrupts operate as follows:

A) Nonvectored Interrupt (I1)

PS and PC are pushed onto the stack. I2 is disabled. The external status register is tested for a power fail. If power fail is true PC is fetched from location "14". If power fail is false PC is fetched from location "2A", and a microm state code is transmitted to clear the line clock (see appendix D).

B) Vectored Interrupt (IØ)

PS and PC are pushed onto the stack. I2 is disabled. An Interrupt Acknowledge is executed, and the device code of the interrupting device is read in and stripped to bits 1-4. PC is fetched from location \*NOTE: Although only a 4 bit device code is currently used, a minor microm change can allow a device code of from 1-15 bits.

"28" and the device code is added to it. The contents of this intermediate location are read in and added to PC to form the final address. Each intermediate location is a table entry that contains the PC relative offset from the start of the device handler routine to itself. The absolute address of the start of the table is in location "28".

## PRIORITY MASK

Associated with the interrupts is a priority interrupt mask. This is a 16 bit mask where each bit position represents a priority level. Each priority level can be assigned to one or more devices. A one in any bit position can represent an interrupt enable or disable for its associated devices as the hardware dictates. The SAVS, RSTS, and MSKO op codes each alter the mask. When the mask is altered it is written into location "2E" for storage. While the mask is on the bus a microm state code is transmitted (see appendix D) to signal the I/O devices that a new mask is being transmitted. Each device can then look at its assigned mask bit while the memory write to location "2E" is taking place. Whether or not the mask feature is actually used by the I/O devices in no way alters the operations of the op codes mentioned above.

## EXTERNAL STATUS REGISTER

As a part of the hardware external to the CPU the External Status Register supplies the CPU, upon demand, with information about the status of certain hardware areas. This register is gated onto the bus when its associated microm state code is present (see appendix D). The format of the register is as follows:

Bit 7 = Power Fail Status

Bit 6 = Bus Error (Time Out) Status

Bit 5 = Parity Error Status

Bit 4 = I2 Interrupt Line Status

Bit 3 = Halt Option Jumper #2

Bit 2 = Halt Option Jumper #1

Bit 1 = Power Up Option Jumper #2

Bit Ø = Power Up Option Jumper #1

Bits 8-15 are don't care. Bits 5-7 are real time error conditions that also generate a system reset (see next section). Bit 4 is the interrupt enable status. The jumpers can be logic units, switches, or hard wired jumpers as the user wishes. The various options associated with the 4 jumpers are discussed later.

#### POWER UP OPTIONS

A system reset indicate one of 4 conditions: power fail, bus error, parity error, or power up. There are 2 levels of power fail possible in this system (see appendix C): minor and major. Only a major power fail generates a system reset. Both types set bit 7 in the External Status Register. The following steps are performed after a system reset.

- Al) Trace and wait flags are reset if on.
- A2) The external Status Register is fetched.

- A3) The Line-clock-clear state code is transmitted.
- A4) I2 is reset.
- A5) If power fail bit is set go to D1.
- A6) If bus error bit is set go to C1.
- A7) If parity error bit is set go to B1.
- A8) Go to D2 otherwise.
- Bl) Push PS and PC onto stack.
- B2) Fetch PC from location "12" and begin execution.
- C1) Push PS and PC onto stack.
- C2) Fetch PC from location "18" and begin execution.
- D1) Wait until power fail status =  $\emptyset$ .
- D2) Send a system reset microm state code.
- D3) Wait 300 cycles.
- D4) Execute power up option 1,2,3 or 4 per jumpers.

For a proper initial power up either bit 7 must be set or bits 5-7 must be reset when the system reset line is released.

The 4 power up options are as follows:

JUMPERS	OPERATION
ØØ	Execute user bootstrap routine.
Ø1	Pick up RØ-R5, SP, PC, and PS from memory locations $\emptyset$ -"l $\emptyset$ ".
1.Ø	Execute selected halt option.
11	Fetch PC from location "16".
	•

#### HALT OPTIONS

When the halt switch (I3) is set during program execution one of 4 halt options is selected. The halt op code\* and power up option #2 also select the halt option specified. The options are as follows:

JUMPERS	OPERATION
ØØ Ø1	Execute user bootstrap routine.
μī	Save RØ-R5,SP,PC and PS in memory locations $\emptyset$ -"l $\emptyset$ ". Wait until I3 = $\emptyset$ , then restore RØ-
lø	R5,SP,PC and PS from memory locations $\emptyset$ -"l $\emptyset$ ". Lock up processor (requires a system reset to clear).
11	Fetch new PC from location "16".

\*NOTE: Conditional. See Chapter 3.

## USER BOOTSTRAP ROUTINE

When the user bootstrap routine is selected as an option the system creates the starting address by placing address "C $\emptyset\emptyset\emptyset$ " in PC and then replacing bits 8-13 with the contents of the 6 bit External Address Register. This register is gated in with a microm status code (see appendix D).

It allows the user 64 different starting addresses in the range "CØØØ" to "FFØØ".

## SYSTEM ERROR TRAPS

With the exception of the major power fail error that is a function of a system reset, all error conditions perform a common routine as outlined below. A non-vectored interrupt and some op codes also use this routine. The numbers in parenthesis refer to notes that follow the table.

- 1) PS is pushed onto the stack
- 2) PC is pushed onto the stack
- 3) PC is fetched from location X where "X" is from the following table
- (1)(2)(3) "12" for bus error PC
- (1)(2)(3) "14" for nonvectored interrupt power fail PC
- (1)(2)(3) "18" for parity error PC
- (1)(2)(3) "lA" for reserved op code error PC
- (1)(2)(3) "1C" for illegal op code format error PC
- (1)(2)(3) "lE" for XCT error PC
- (1)(2) "2Ø" for XCT trace PC
- (1)(2)(3) "2A" for nonvectored interrupt PC
- (1)(2) "2C" for BPT PC
- NOTE 1: wait flag reset if on
- NOTE 2: trace flag reset if on
- NOTE 3: interrupt enable (I2) reset if on

The meaning of the wait and trace flags is discussed in chapter 3. Note that the nonvectored interrupt power fail PC is a minor power fail condition, not a major one. See appendix C for full detail on how to include both major and minor power fail conditions in the hardware.

## RESERVED CORE LOCATIONS

The following is a complete list of memory locations that are reserved for specific system functions or options. Byte addresses are given.

LOCATIONS	RESERVED FUNCTION
Ø - "11"	RØ - R5, SP, PC and PS for power up/halt options
"12" - "13"	bus error PC
"14 - "15"	nonvectored interrupt power fail PC
"16" - "17"	power up/halt option power restore PC
"18" - "19"	parity error PC
"1A" -"1B"	reserved op code PC
"1C" - "1D"	illegal op code format PC
"1E" - "1F"	XCT error PC
"2Ø" - "21"	XCT trace PC
"22" - "23"	SVCA table address
"24" - "25"	SVCB PC
"26" - "27"	SVCC PC
"28"- "29"	vectored interrupt (IØ) table address
"2A" - "2B"	nonvectored interrupt (I1) PC
"2C" - "2D"	BPT PC
"2E" - 2F"	I/O priority interrupt mask
"3Ø" – "3F"	reserved for floating point option
·	5

## CHAPTER 3 - OP CODES

This chapter is divided into a number of sections, each representing one class of op codes. At the beginning of each section there is a detailed description of the format for that class. A list of op codes and their base numeric values, less arguments, is also included. A detailed description of each op code in the class then follows.

## FORMAT 1 OP CODES

Single word - no arguments

15	12	11 8	7	4	3 0
	Ø	Ø		Ø	OPC

There are 16 op codes in this class representing op codes "ØØØØ" to "ØØØF". Each is a one word op code with no arguments with the exception of the SAVS op code which is a two word op code. Word two of the SAVS op code is the I/O priority interrupt mask. The op codes and their mnemonics are:

BASE OP CODE	MNEMONIC
ØØØØ	NOP
ØØØ1	RESET
ØØØ2	IEN
ØØØ3	IDS
ØØØ4	HALT
øøø5	XCT
øøø6	BPT
ØØØ7	WFI
øøø8	RSVC
øøø9	RRTT
ØØØA	SAVE
ØØØB	SAVS
øøøc	REST
ØØØD	RRTN
ØØØE	RSTS
ØØØF	RTT
NOP	NO OPERATION
FORMAT:	NOP
FUNCTION:	No operations are performed
INDICATORS:	Unchanged
RESET	I/O RESET
FORMAT:	RESET
FUNCTION:	An I/O reset pulse is transmitted
INDICATORS:	Unchanged

IEN INTERRUPT ENABLE

FORMAT:

IEN

FUNCTION:

The interrupt enable (I2) flag is set. Allows one more instruction to execute before inter-

rupts are recognized.

INDICATORS:

Unchanged

IDS

INTERRUPT DISABLE

FORMAT:

IDS

FUNCTION:

The interrupt enable (I2) flag is reset. This instruction can honor interrupts, but the I2 bit in the PS that is stored on the stack

is reset if an interrupt occurs.\*

INDICATORS:

Unchanged

\*NOTE: On some machines I2 will be set or reset during the IEN or IDS. If so the change will be valid immediately, not one op code later.

HALT HALT

FORMAT:

דאדת

FUNCTION:

Tests the status of the Power Fail bit in the external status register. If the bit is set it is assumed that the HALT occured in a power fail routine, and the following operations occur:

- 1) The interrupt enable (I2) flag is reset
- 2) The CPU waits until the Power Fail bit is reset
- 3) PC is fetched from location "16", and program execution begins at this new location

If the power fail bit is reset then the CPU waits until the halt switch (I3) is set. At that time the selected halt option (see chapter 2) is executed.

The interrupt enable flag is also reset.

INDICATORS:

Unchanged

XCT

## EXECUTE SINGLE INSTRUCTION

FORMAT:

OPERATION:

XCT

 $PC \leftarrow @SP, SP \uparrow$  $PS \leftarrow @SP, SP \uparrow$ 

Trace flag set, execute op code

 $\forall$  SP, @SP  $\leftarrow$  PS  $\forall$  SP, @SP  $\leftarrow$  PC Trace flag reset

PC ←(loc "2Ø") if no error PC ←(loc "lE") if error

FUNCTION:

PC and PS are popped from the stack, but I2 is not altered. The trace flag, which disables all interrupts except I3, is set. The op code is executed. PS and PC are pushed back onto the stack, and PC is fetched from location "20". The trace flag is reset. If the program tries to execute a HALT, XCT, BPT, or WFI the attempt is aborted, PS and PC are

pushed onto the stack, and PC is fetched from location "lE" instead.

INDICATORS:	Depends upon executed op code
BPT	BREAKPOINT TRAP
FORMAT:	врт
OPERATION:	↓SP, @SP ←PS
	↓SP, @SP ←PC
	PC ← (loc "2C")
FUNCTION:	PS and PC are pushed onto the stack. PC is fetched from location "2C"
INDICATORS:	Unchanged
WFI	WAIT FOR INTERRUPT
	MINT
FORMAT:	WFI
FUNCTION:	The CPU loops internally without accessing the data bus until an interrupt occurs. Program execution continues with the op code that follows
	the WFI after the interrupt has been serviced.
	The interrupt enable flag is also set.
TNIDTCA MODC.	Unchanged
INDICATORS:	Olichangeu
SAVE	SAVE REGISTERS
FORMAT:	SAVE
OPERATION:	↓ SP , @SP + R5
	↓SP, @SP ← R4
	$\downarrow \text{SP}$ , $\text{@SP} \leftarrow \text{R3}$
	♦SP, @SP + R2
	↓SP, @SP ← R1
	↓SP, @SP ← RØ
FUNCTION:	Registers R5 to RØ are pushed onto the stack.
INDICATORS:	Unchanged.
SAVS	SAVE STATUS
TONKS M	CALLC MACK
FORMAT:	SAVS MASK
OPERATION:	SAVE √SP, @SP ← (loc "2E")
	(loc "2E") ← (loc "2E") ▼ mask
	MSKO
	MSRO IEN
TODA 5	
FORMAT:	Registers R5 to RØ and the priority mask in location "2E" are pushed onto the stack. The old and new masks
	are ORED together and placed in location "2E".
	A mask out state code (see appendix D) is transmitted
	and the interrupt enable (I2) flag is set.
INDICATORS:	Unchanged
INDICATORS:	onenanged
REST	RESTORE REGISTERS
FORMAT:	REST
OPERATION:	$R\emptyset \leftarrow QSP$ , $SP \uparrow$
	Rl ← @SP, SP ↑

R2 + @SP, SP \*

 $R3 \leftarrow @SP, SP \uparrow$  $R4 \leftarrow @SP, SP \uparrow$  $R5 \leftarrow @SP, SP \uparrow$ 

FUNCTION: Registers Rg to  $\ensuremath{\text{R}}\xspace5$  are popped from the stack, INDICATORS: Unchanged

RTT	RETURN FROM TRAP
FORMAT:	RTT
OPERATION:	PC ←@SP, SP↑
	PS ←@SP, SP↑
FUNCTION:	PC and PS are popped from stack
INDICATORS:	N = Set per PS bit 3
	Z = Set per PS bit 2
	V = Set per PS bit 1
	C = Set per PS bit Ø
RRTN	RESTORE AND RETURN FROM SUBROUTINE
FORMAT:	RRTN
OPERATION:	REST
or Brutton.	PC ← @SP, SP↑
FUNCTION:	Registers RØ to R5 and PC are popped
	from the stack
INDICATORS:	Unchanged
RRTT	RESTORE AND RETURN FROM TRAP
FORMAT:	RRTT
OPERATION:	REST
or Bittizon.	RTT
FUNCTION:	Registers RØ to R5, PC and PS are popped
	from the stack.
INDICATORS:	Set per PS bits ∅ - 3
RSTS	RESTORE STATUS
FORMAT:	RSTS
OPERATION:	(LOC "2E") ← @SP, SP↑
	MSKO
	REST
	RTT
FUNCTION:	The priority mask is popped from the stack and
	restored to locaton "2E". A MASK OUT state code
	(See Appendix D) is transmitted. Registers RØ
	to B, PC and PS are popped from the stack.
INDICATORS:	Set per PS bits Ø - 3
RSVC	RETURN FROM SUPERVISOR CALL (B or C)
FORMAT:	RSVC
OPERATION:	REST
	SP
	RTT

FUNCTION:

INDICATORS:

Registers  $R\beta$  to R5 , PC and PS are popped from the stack with the saved SP bypassed.

Set per PS bits  $\emptyset$  - 3

## FORMAT 2 OP CODES

## SINGLE WORD - 3 BIT REGISTER ARGUMENT

15 12	11 8	7 3	2 0
ø	Ø	OPC	REG

There are 4 op codes in this class representing op codes " $\emptyset\emptyset$ 1 $\emptyset$ " to " $\emptyset\emptyset$ 2F". Each is a one word op code with a single 3 - bit register argument. The op codes and their mnemonics are:

BASE OP CODE	MNEMONIC
ØØ1Ø ØØ18 ØØ2Ø ØØ28	IAK RTN MSKO PRTN
IAK	INTERRUPT ACKNOWLEDGE
FORMAT: FUNCTION:	IAK REG An interrupt acknowledge (READ and IACK) is executed, and the 16 bit code that is returned is placed in REG unmodified. Used with the nnnvectored interrupt when the user does not wish to use the vectored format.
INDICATORS:	Unchanged
RTN	RETURN FROM SUBROUTINE
FORMAT: OPERATION:	RTN REG PC ← REG REG ← @SP,SP↑
FUNCTION:	The linkage register is placed in PC and the saved linkage register is popped from the stack. The register used must be the same one that was used for the subroutine call.
INDICATORS:	Unchanged
MSKO	MASK OUT
FORMAT: OPERATION:	MSKO REG (LOC "2E" ) ← REG MSKO
FUNCTION:	The contents of REG are written into location "2E" and a MASK OUT state code (see appendix D) is transmitted.
INDICATORS:	Unchanged
PRTN	POP STACK AND RETURN
FORMAT: OPERATION:	PRTN REG TMP ← @SP SP ← SP+(TMP*2) RTN REG

FUNCTION:

Twice the value of the top word on the stack is added to SP, and a standard RTN call is then executed.

INDICATORS:

Unchanged

## FORMAT 3 OP CODES

## SINGLE WORD - 4 BIT NUMERIC ARGUMENT

15	12	11	8	7	4	3	Ø
Ø		1	Ø		OPC	Al	RG

There is only one op code in this class representing op codes " $\emptyset\emptyset3\emptyset$ " to " $\emptyset\emptyset3F$ ". It is a one word op code with a 4-bit numeric argument.

BASE OP CODE	MNEMONIC
ØØ3Ø	LCC
LCC	LOAD CONDITION CODES
FORMAT:	LCC ARG
FUNCTION:	The 4 indicators are loaded from bits $\emptyset$ -3 of the op code as specified.
INDICATORS:	<pre>N = Set per bit 3 of op code Z = Set per bit 2 of op code V = Set per bit 1 of op code C = Set per bit Ø of op code</pre>

## FORMAT 4 OP CODES

## SINGLE WORD - 6 BIT NUMERIC ARGUMENT

15 12	11 8	7 6	5 Ø
Ø	18	OPC	ARG

There are 3 op codes in this class representing op codes " $\emptyset\emptyset4\emptyset$ " to " $\emptyset\emptysetFF$ ". All 3 are supervisor calls . All 3 are one word op codes with a 6-bit numeric argument.

BASE OP CODE	MNEMONIC
ØØ4Ø ØØ8Ø ØØCØ	SVCA SVCB SVCC
SVCA	SUPERVISOR CALL "A"
FORMAT: OPERATION:	SVCA ARG ↓SP, @SP ← PS; ↓ SP, @SP← PC PC ← (LOC "22") + (ARG *2) PC ← PC + @PC
FUNCTION:	PS and PC are pushed onto the stack. The contents of location "22" plus twice the value of the argument (which is always positive) is placed in PC to get the table address. The contents of the table address is added to PC to get the final destination address. Each table entry is the relative offset from the start of the desired
INDICATORS:	routine to itself. Unchanged
SVCB SVCC	SUPERVISOR CALL "B" SUPERVISOR CALL "C"
FORMAT:	SVCB ARG SVCC ARG
OPERATION:	TMPA ← SP ↓SP, @SP ← PS ↓SP, @SP ← PC
	TMPB ←SP →SP, @SP ← TMPA SAVE R1 ← TMPB
	R5 ← ARG*2 PC ← (LOC "24") if SVCB PC ← (LOC "26") if SVCC
FUNCTION:	PS and PC are pushed onto the stack. The value

of SP at the start of op code execution is the pushed followed by registers R5 to RØ. The address of the saved PC is placed in R1, and twice the value of the 6-bit positive argument is placed in R5.

PC is loaded from location "24" for SVCB or "26" for SVCC. Unchanged.

INDICATORS:

## FORMAT 5 OP CODES

## SINGLE WORD - 8 BIT SIGNED NUMERIC ARGUMENT

15		8	7	Ø	
	OPC			DISPLACEMENT	

There are 15 op codes in this class representing op codes "Ø1ØØ" to "Ø7FF" and "8ØØØ" to "87FF". All are branches with a signed 8 bit displacement that represents the word offset from PC (which points to the op code that follows) to the desired branch location. The op codes consist on one unconditional branch, 8 signed conditional branches, and 6 unsigned conditional branches. No op code in this class modifies any of the indicator flags. Maximum branch range is +128, -127 words from the branch op code.

BASE OP CODE	MNEMONIC
ØlØØ	BR
Ø2ØØ	BNE
Ø3ØØ	BEQ
Ø4ØØ	BGE
Ø5ØØ	BLT
ø6øø	BGT
Ø7ØØ	BLE
8ØØØ	BPL
81ØØ	BMI
82ØØ	BHI
83ØØ	BLOS
84ØØ	BVC
85ØØ	BVS
86ØØ	BCC, BHIS
87ØØ	BCS, BLO
ששוט	
BR	BRANCH UNCONDITIONALLY
FORMAT:	BR DEST
OPERATION:	PC ← PC+ (DISP *2)
FUNCTION:	Twice the value of the signed displacement
	is added to PC.
	SIGNED BRANCHES
	BRANCH IF NOT EQUAL TO ZERO
BNE	BRANCH II NOT EXCEL TO
	BNE DEST
FORMAT:	IF $Z = \emptyset$ , $PC \leftarrow PC + (DISP *2)$
OPERATION:	
BEQ	BRANCH IF EQUAL TO ZERO
BEQ	
FORMAT:	BEO DEST
OPERATION:	$\overline{IFZ} = 1$ , PC $\leftarrow$ PC + (DISP $*2$ )
OF BIGHTON.	
BGE	BRANCH IF GREATER THAN OR EQUAL TO ZERO
FORMAT:	BGE DEST
OPERATION:	IF $N\nabla V = \emptyset$ , PC $\leftarrow$ PC + (DISP *2)

BLT	BRANCH IF LESS THAN ZERO
FORMAT: OPERATION:	BLT DEST IF N∀V = 1, PC ← PC + (DISP *2)
BGT	BRANCH IF GREATER THAN ZERO
FORMAT: OPERATION:	BGT DEST IF Z ∇(NΨV) = Ø, PC + PC + (DISP *2)
BLE	BRANCH IF LESS THAN OR EQUAL TO ZERO
FORMAT: OPERATION:	BLE DEST IF $Z\nabla(N\nabla V) = 1$ , PC $\leftarrow$ PC + (DISP *2)
BPL	BRANCH IF PLUS
FORMAT: OPERATION:	BPL DEST IF N = $\emptyset$ , PC $\leftarrow$ PC + (DISP *2)
BMI	BRANCH IF MINUS
FORMAT: OPERATION:	BMI DEST IF N = 1, PC $\leftarrow$ PC + (DISP *2)
	UNSIGNED BRANCHES
ВНІ	BRANCH IF HIGHER
FORMAT: OPERATION:	BHI DEST IF $C\nabla Z = \emptyset$ , $PC \leftarrow PC + (DISP *2)$
BLOS	BRANCH IF LOWER OR SAME
FORMAT: OPERATION:	BLOS DEST IF $C\nabla Z = 1$ , $PC \leftarrow PC + (DISP *2)$
BVC	BRANCH IF OVERFLOW CLEAR
FORMAT: OPERATION:	BVC DEST IF $V = \emptyset$ , PC $\leftarrow$ PC + (DISP *2)
BVS	BRANCH IF OVERFLOW SET
FORMAT: OPERATION:	BVS DEST IF V = 1, PC \( \text{PC} + (DISP *2)
BCC	BRANCH IF CARRY CLEAR
BHIS	BRANCH IF HIGHER OR SAME
FORMAT:	BCC DEST
OPERATION:	BHIS DEST IF $C = \emptyset$ , $PC \leftarrow PC + (DISP *2)$

BCS	BRANCH IF CARRY SET
BLO	BRANCH IF LOWER
FORMAT:	BCS DEST BLO DEST
OPERATION:	IF $C = 1$ , $PC \leftarrow PC + (DISP *2)$

## FORMAT 6 OP CODES

SINGLE WORD - SINGLE OPS - SPLIT FIELD - DMØ ONLY

15	9	8	6	5	4	3	0
OPC BAS	E	R	EG		OPC	CC	UNT

There are 12 op codes in this class representing op codes " $\emptyset 8 \emptyset \emptyset$ " to " $\emptyset 9 F F$ ", " $88 \emptyset \emptyset$ " to "89 F F", and " $8E \emptyset \emptyset$ " to "8F F F". There are 4 immediate mode op codes with a register as a destination, 4 multiple count single register shifts, and 4 multiple count double register shifts. In all op codes the actual count (or number in the case of the immediates) is the value of bits  $\emptyset$  - 3 plus one. Count is always a positive number in the range 1 - " $1\emptyset$ ", but it is stored in the op code as  $\emptyset$  - "F". All of these op codes are one word op codes with the op codes themselves split between bits 9-15 and 4-5.

In the case of the double shifts the 32 bit number (REG+1) : (REG) is the operand. If REG = PC then (REG+1) =  $R\emptyset$ .

BASE OP CODE	MNEMONIC
ø8øø	ADDI
Ø81Ø	SUBI
Ø82Ø	BICI
Ø83Ø	MOVI
88ØØ	SSRR
881Ø	SSLR
882Ø	SSRA
883Ø	SSLA
8EØØ	SDRR
8E1Ø	SDLR
8E2Ø	SDRA
8E3Ø	SDLA
ADDI	ADD IMMEDIATE
FORMAT;	ADDI NUMBER, REG
OPERATION:	REG ← REG + COUNT + 1
FUNCTION:	The stored number plus one is added to the
T	destination register.
INDICATORS:	N = Set if bit 15 of the result is set
	$Z = Set if the result = \emptyset$
	V = Set if arithmetic overflow occurs; i.e. set
	if both operands were positive and the sign of
	the result is negative
	C = Set if a carry was generated from bit 15
	of the result
SUBI	SUBTRACT IMMEDIATE
FORMAT:	SUBI NUMBER, REG
OPERATION:	REG ← REG - (COUNT +1)
FUNCTION:	The stored number plus one is subtracted from
··, <del></del>	the destination register
	•

INDICATORS:

N = Set if bit 15 of the result is set

 $Z = Set if the result = \emptyset$ 

V = Set if arithmetic underflow occurs; i.e. set if the operands were of opposite signs and

the sign of the result is positive

C = Set if a borrow was generate from bit 15

of the result

BICI

BIT CLEAR IMMEDIATE

FORMAT:

OPERATION:

BICI NUMBER, REG REG  $\leftarrow$  REG  $\triangle$  (COUNT + 1)

FUNCTION:

The stored number plus one is one's complemented

and ANDED to the destination register

INDICATORS:

N = Set if bit 15 of the result is set

 $Z = Set if the result = \emptyset$ 

V = Reset

C = Unchanged

IVOM

MOVE IMMEDIATE

FORMAT:

OPERATION:

IVOM NUMBER, REG

REG ← COUNT + 1

FUNCTION:

The stored number plus one is placed in

the destination register

INDICATORS:

N = Reset Z = Reset

V = Reset

C = Unchanged

SSRR

SHIFT SINGLE RIGHT ROTATE

FORMAT:

FUNCTION:

SSRR REG, COUNT

A 17-bit right rotate is done stored count+1 times on REG: C-Flag. The C-Flag is shifted into

bit 15 of REG, and the C-Flag gets the last bit

shifted out of REG bit Ø.

INDICATORS:

N = Set if bit 7 of REG is set

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags C = Set to the value of the last bit shifted

out of REG bit Ø

SSLR

SHIFT SINGLE LEFT ROUTINE

FORMAT:

SSLR REG, COUNT

FUNCTION:

INDICATORS:

A 17-bit left rotate is done stored count+1 times on C-Flag: REG . The C-Flag is shifted into bit  $\emptyset$  of REG and the C-Flag gets the

last bit shifted out of REG bit 15.

N = Set if bit 15 of REG is set

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted

out of REG bit 15.

## SSRA SHIFT SINGLE RIGHT ARITHMETIC

FORMAT: SSRA REG, COUNT

FUNCTION: A 17-bit right arithmetic shift is done stored count+1 times on REG: C-Flag. Bit

15 of REG is replicated. The C-Flag gets the last bit shifted out of REG bit Ø. Bits shifted

out of the C-Flag are lost.

INDICATORS: N = Set if bit 7 of REG is set

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted

out of REG bit Ø

## SSLA SHIFT SINGLE LEFT ARITHMETIC

FORMAT: SSLA REG, COUNT

FUNCTION: A 17-bit left arithmetic shift is done stored

count+1 times on C-Flag:REG. Zeros are shifted into REG bit  $\emptyset$ , and the C-FLAG gets the last bit shifted out of REG bit 15. Bits shifted out of the

C-Flag are lost

INDICATORS:  $N = Set \ if \ REG \ bit \ 15 \ is \ set$ 

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted

out of REG bit 15

#### SDRR SHIFT DOUBLE RIGHT ROTATE

FORMAT: SDRR REG, COUNT

FUNCTION: REG+1:REG:C-Flag is rotate right stored

count+1 times. The C-Flag is shifted into REG+1 bit 15, REG+1 bit Ø is shifted into

REG bit 15, and REG bit Ø is shifted into the C-Flag.

INDICATORS: N = Set if bit 7 of REG is set

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted

out of REG bit Ø

## SDLR SHIFT DOUBLE LEFT ROTATE

FORMAT: SDLR REG, COUNT

FUNCTION: A 33 bit left rotate is done stored count+1

times on C-Flag:REG+1:REG. The C-Flag is shifted into REG bit Ø, REG bit 15 is shifted into REG+1 bit Ø, and REG+1 bit 15 is shifted

into the C-Flag

INDICATORS: N = Set if REG+1 bit 15 is set

 $Z = Set if REG+1 = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted

out of REG+1 bit 15.

DRA	SHIFT	DOUBLE	RIGHT	ARITHMETIC

FORMAT:

SDRA REG, COUNT

FUNCTION:

A right arithmetic shift is done stored

count+1 times on REG+1:REG:C-Flag.

Bit 15 of REG+1 is replicated. Bit  $\emptyset$  of REG+1 is shifted to bit 15 of REG. Bit  $\emptyset$  of REG is shifted to the C-Flag. Bits

shifted out of the C-Flag are lost.

INDICATORS:

N = Set if bit 7 of REG is set

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit

shifted out of REG bit Ø

#### SDLA

## SHIFT DOUBLE LEFT ARITHMETIC

FORMAT: FUNCTION:

SDLA REG, COUNT

A left arithmetic shift is done stored

count+1 times on C-Flag: REG+1: REG.

Zeros are shifted into REG bit  $\emptyset$ , REG bit 15 is shifted to REG+1 bit  $\emptyset$ . REG+1 bit 15 is shifted to the C-Flag. Bits shifted out of the C-Flag are lost.

INDICATORS:

N = Set if REG+l bit 15 is set

 $Z = Set if REG+l = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted

out of REG+1 bit 15

#### FORMAT 7 OP CODES

SINGLE OPS - ONE OR TWO WORDS - DMØ TO DM7

1.	5	6	5	3	2	0
Г	OPC		N	ODE		REG

There are 32 op codes in this class representing op codes "@A@@" to "@DFF" and "8A@@" to "8DFF". All addressing modes from @ to 7 are available with all registers available as index registers (see chapter two). A one word op code is generated for addressing modes @ to 5. A two word op code is generated for addressing modes 6 and 7 with the offset value in word two. For DM6 and DM7 with PC as the index register PC is added to the offset from word two after the offset is fetched from memory. The offset is therefore relative to a PC that points to the op code that follows (i.e. current op code + 4). Codes "8A@@" to "8CC@" are BYTE ops.

BASE OP CODE	MNEMONIC	BASE OP CODE	MNEMONIC	
ØAØØ	ROR	8aøø	RORB	
ØA4Ø	ROL	8A4Ø	ROLB	
ØA8Ø	$ extbf{T} extbf{S} extbf{T}$	8A8Ø	TSTB	
ØACØ	ASL	8ACØ	ASLB	
øbøø	SET	8B <b>Ø</b> Ø	SETB	
ØB4Ø	CLR	8B4Ø	CLRB	
ØB8Ø	ASR	8B8Ø	ASRB	
ØBCØ	SWAB	8BCØ	SWAD	
ØCØØ	COM	8CØØ	COMB	
ØC4Ø	NEG	8C4Ø	NEGB	
ØC8Ø	INC	8C8Ø	INCB	
ØCCØ	DEC	8CCØ	DECB	
ØDØØ	IW2	8DØØ	LSTS	
ØD4Ø	SXT	SD4Ø	SSTS	
ØD8Ø	TCALL	8D8Ø	ADC	
ØDCØ	TJMP	8DCØ	SBC	
م د ح م				

## WORD OPS

ROR	ROTATE RIGHT
FORMAT:	ROR DST
FUNCTION:	A 1-bit right rotate is done on (DST):C-Flag
	The C-Flag is shifted into (DST) bit 15, and (DST)
	bit Ø is shifted into the C-flag.
INDICATORS:	N = Set if bit 7 of (DST) is set
INDICATORS.	$Z = Set if (DST) = \emptyset$
	V = Set to exclusive or of N and C flags
	C = Set to the value of the bit shifted out of (DST)
ROL	ROTATE LEFT
FORMAT:	ROL DST
FUNCTION:	A 1-bit left rotate is done on C-Flag: (DST). The

C-Flag is shifted into (DST) bit Ø, and (DST)

bit 15 is shifted into the C-Flag.

INDICATORS: N = Set if bit 15 of (DST) is set

 $Z = Set if (DST) = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the bit shifted out of (DST)

TEST WORD

TSTFORMAT: OPERATION: (DST) A (DST)

The indicators are set to reflect the destination FUNCTION:

DST

operand status.

INDICATORS: N = Set if (DST) bit 15 is set

 $Z = Set if (DST) = \emptyset$ 

V = Reset C = Unchanged

ASL ARITHMETIC SHIFT LEFT

FORMAT: ASL DST

A 1-bit left arithmetic shift is done on (DST). A FUNCTION:

zero is shifted into (DST) bit  $\emptyset$ , and (DST) bit 15

is shifted into the C-Flag.

N = Set if (DST) bit 15 is set INDICATORS:

 $Z = Set if (DST) = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the bit shifted out of (DST)

SET TO ONES SET

FORMAT: SET DST (DST) ← "FFFF" OPERATION:

FUNCTION: The destination operand is set to all ones

N = SetINDICATORS: Z = Reset

> V = Reset C = Unchanged

CLR CLEAR TO ZEROS

FORMAT: CLR DST OPERATION:  $(DST) \leftarrow \emptyset$ 

The destination operand is cleared to all zeros FUNCTION:

INDICATORS: N = ResetZ = SetV = Reset

C = Unchanged if DMØ. Reset if DMl-DM7.

ARITHMETIC SHIFT RIGHT

FORMAT: ASR DST

FUNCTION: A 1-bit right arithmetic shift is done on (DST). Bit

15 of (DST) is replicated. Bit Ø of (DST) is shifted

into the C-Flag.

```
N = Set if (DST) bit 7 is set
INDICATORS:
                    Z = Set if (DST) = \emptyset
                    V = Set to exclusive or of N and C flags
                    C = Set to the value of the bit shifted out of (DST)
                    SWAP BYTES
SWAB
FORMAT:
                    SWAB
                            DST
                    (DST) 15-8 \stackrel{?}{\leftarrow} (DST) 7-\emptyset
OPERATION:
                    The upper and lower bytes of (DST) are exhanged.
FUNCTION:
                    N = Set if (DST) bit 7 is set
INDICATORS:
                    Z = Set if (DST) lower byte = \emptyset
                    V = Reset
                    C = Unchanged
                    COMPLEMENT
COM
FORMAT:
                    COM
                          DST
                    (DST) \leftarrow (DST)
OPERATION
                    The destination operand is one's complemented.
FUNCTION:
                    N = Set if (DST) bit 15 is set
INDICATORS:
                    Z = Set if (DST) = \emptyset
                    V = Reset
                    C = Set
                    NEGATE
NEG
                           DST
                    NEG
FORMAT:
OPERATION:
                    (DST) \leftarrow -(DST)
                    The destination operand is two's complemented.
FUNCTION:
                    N = Set if (DST) bit 15 is set
INDICATORS:
                    Z = Set if (DST) = \emptyset
                    V = Set if (DST) = "8000"
                    C = Reset if (DST) = \emptyset
INC
                    INCREMENT
                    INC
                            DST
FORMAT:
OPERATION:
                    (DST) \leftarrow (DST) + 1
                    The destination operand is incremented by one.
FUNCTION:
                    N = Set if (DST) bit 15 is set
INDICATORS:
                    Z = Set if (DST) = \emptyset
                    V = Set if (DST) = "8000"
                    C = Set if a carry is generated from (DST) bit 15
                    DECREMENT
DEC
FORMAT:
                    DEC
                            DST
                    (DST) \leftarrow (DST) - 1
OPERATION:
                    The destination operand is decremented by one.
FUNCTION:
```

C = Set if a borrow is generated from (DST) bit 15

N = Set if (DST) bit 15 is set

 $Z = Set if (DST) = \emptyset$ V = Set if (DST) = "7FFF"

INDICATORS:

INCREMENT WORD BY TWO IW2 DST FORMAT: IW2  $(DST) \leftarrow (DST) + 2$ OPERATION: The destination operand is incremented by two. FUNCTION: N = Set if (DST) bit 15 is setINDICATORS:  $Z = Set if (DST) = \emptyset$ V = Set if (DST) = "8000" or "8001"C = Set if a carry is generated from (DST) bit 15 SIGN EXTEND SXT SXT DST FORMAT:  $(DST) \leftarrow \emptyset$  $IF N = \emptyset,$ OPERATION: (DST) ← "FFFF" IF N = 1,The N-Flag status is replicated in the destination operand FUNCTION: Unchanged INDICATORS: TABLED SUBROUTINE CALL TCALL DST TCALL FORMAT: ↓ SP, @SP ← PC OPERATION:  $PC \leftarrow PC + (DST)$ PC ← PC + @PC PC, which points to the op code that follows, is pushed FUNCTION: onto the stack. The destination operand is added to PC. The contents of this intermediate table address is also added to PC to get the final destination address. Note that at least one op code must exist between the TCALL and the table for a subroutine return. Unchanged INDICATORS: TABLED JUMP TJMP FORMAT: TJMP DST  $PC \leftarrow PC + (DST)$ OPERATION:  $PC \leftarrow PC + @PC$ The destination operand is added to PC, and the contents FUNCTION: of this intermediate location is also added to PC to get the final destination address. Unchanged INDICATORS: LOAD PROCESSOR STATUS LSTS LSTS DST FORMAT: The four indicators and the interrupt enable (I2) FUNCTION: are loaded from the destination operand. Set to the status of (DST) bits  $\emptyset$  - 3 INDICATORS: STORE PROCESSOR STATUS SSTS SSTS FORMAT:

Unchanged

FUNCTION:

INDICATORS:

The processor status word is formed and stored in (DST).

ADD CARRY

FORMAT:

ADC DST

OPERATION:

 $(DST) \leftarrow (DST) + C-flag$ 

FUNCTION:

The carry flag is added to the destination operand

INDICATORS:

N= Set if (DST) bit 15 is set

 $Z = Set if (DST) = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set if a carry is generated from (DST) bit 15

SBC

SUBTRACT CARRY

FORMAT:

SBC DST

OPERATION:

 $(DST) \leftarrow (DST) - C-Flag$ 

FUNCTION:

The Carry flag is subtracted from the destination operand

INDICATORS:

N = Set if (DST) bit 15 is set

 $Z = Set if (DST) = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set if a borrow is generated from (DST) bit 15

#### BYTE OPS

For DMØ addressing only the lower byte of the destination register is affected by a byte op code. For DM1-DM7 addressing only the specified memory byte is affected by a byte op. For even memory addresses the lower byte is altered, and for ddd memory addresses the upper byte is altered.

RURR

## ROTATE RIGHT BYTE

FORMAT:

RORB DST

FUNCTION:

A 1-bit right rotate is done on (DST)<sub>B:</sub>C-Flag. Bit  $\emptyset$  of (DST)<sub>B</sub> is shifted into the C-Flag, and the C-Flag

is shifted into  $(DST)_B$  bit 7.

INDICATORS:

 $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the bit shifted out of (DST)<sub>B</sub>bit  $\emptyset$ 

#### ROLB

## ROTATE LEFT BYTE

FORMAT:

ROLB DST

FUNCTION:

A 1-bit left rotate is done on C-flag :(DST) $_{\rm B}$ . Bit 7 of (DST)B is shifted into the C-flag, and the C-flag

is shifted into (DST) B bit Ø

INDICATORS:

 $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the bit shifted out of  $(DST)_B$  bit 7

TEST BYTE

TSTB

**EORMAT:** 

 $(DST)_B \Delta (DST)_B$ 

OPERATION:

FUNCTION:

The destination operand status sets the indicators.

INDICATORS:

 $N = Set if (DST)_{R} bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ 

V = Reset C = Unchanged

ASLB

ARITHMETIC SHIFT LEFT BYTE

FORMAT:

ASLB DST

FUNCTION:

A 1-bit left arithmetic shift is done on C-Flag: (DST) B A zero is shifted into (DST)<sub>B</sub> bit  $\emptyset$ , and (DST)<sub>B</sub> bit 7 is

shifted into the C-flag.

INDICATORS:

 $N = \text{set if } (DST)_B \text{ bit 7 is set}$ 

 $z = Set if (DST)_B = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the bit shifted out of  $(DST)_B$  bit 7

SETB

SET BYTE TO ONES

FORMAT:

SETB DST

OPERATION:

 $(DST)_R \leftarrow "FF"$ 

FUNCTION:

The destination byte operand is set to all ones

INDICATORS:

N = SetZ = Reset V = Reset C = Unchanged

CLRB

CLEAR BYTE TO ZEROS

FORMAT:

CLRB DST

OPERATION:

 $(DST)_R \leftarrow \emptyset$ 

FUNCTION:

The destination byte operand is cleared to all zeros.

INDICATORS:

N = Reset Z = Set

V = Reset

C = Reset

ASRB

ARITHMETIC SHIFT RIGHT BYTE

FORMAT:

ASRB DST

FUNCTION:

INDICATORS:

A 1-bit right arithmetic shift is done on (DST)<sub>R</sub>:

C-flag. Bit 7 of (DST)  $_{\rm B}$  is replicated. Bit  $\emptyset$  of (DST)  $_{\rm B}$  is shifted into the C-flag.

 $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_{R}^{D} = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the bit shifted out of (DST)<sub>R</sub> bit  $\emptyset$ 

SWAD

SWAP DIGITS

FORMAT:

SWAD

FUNCTION:

The two hex digits in the destination byte operand

are exchanged with each other,

INDICATORS:

 $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ 

 $V = Set if (DST)_B bit 7 is set$ 

C = Reset

#### COMPLEMENT BYTE COMB

FORMAT: COMB DST OPERATION:  $(DST)_B \leftarrow \overline{(DST)_B}$ 

The destination byte operand is one's complemented FUNCTION:

INDICATORS:  $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ 

V = Reset C = Set

#### NEGATE BYTE NEGB

FORMAT: NEGB DST (DST)<sub>B</sub> - (DST)<sub>B</sub> OPERATION:

FUNCTION: The destination byte operand is two's complemented

INDICATORS:  $N = Set if (DST)_B bit 7 is set$ 

 $z = Set if (DST)_B = \emptyset$  $V = Set if (DST)_B = "8000"$  $C = Reset if (DST)_B = \emptyset$ 

#### INCB INCREMENT BYTE

FORMAT: INCB DST OPERATION:  $(DST)_B \leftarrow (DST)_B + 1$ 

FUNCTION: The destination byte operand is incremented by one

 $N = Set if (DST)_B is set$ INDICATORS:

Z = Set if (DST)<sub>B</sub> =  $\emptyset$ V = Set if (DST)<sub>B</sub> = "8 $\emptyset$  $\emptyset$  $\emptyset$ "

 $C = Set if a carry is generated from (DST)_R bit 7$ 

#### **DECB** DECREMENT BYTE

FORMAT: DECB DST

**OPERATION:**  $(DST)_B \leftarrow (DST)_B - 1$ 

FUNCTION: The destination byte operand is decremented by one

INDICATORS:  $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ V = Set if (DST)\_B = "7FFF"

C = Set if a borrow is generated from (DST)<sub>B</sub> bit 7

#### FORMAT 8 OP CODES\*

DOUBLE OPS - SINGLE WORD - SMØ AND DMØ ONLY

15	6	5		3	2	ø
OPC		S	REG		D	REG

There are 8 op codes in this class representing op codes " $\emptyset E \emptyset \emptyset$ " to " $\emptyset F F F$ ". Only addressing mode  $\emptyset$  is allowed for both the source and destination. All are one word op codes, and all are block move instructions. The last 4 can be used as pseudo DMA ops in some hardware configurations. In all cases the source register contains the address of the first word or byte of memory to be moved, and the destination register contains the address of the first word or byte of memory to receive the data being moved. The number of words or bytes being moved is contained in R $\emptyset$ . The count ranges from 1-65536 ( $\emptyset$  = 65536) words or bytes. The count in R $\emptyset$  is an unsigned positive integer. None of the indicators are altered by these op codes.

Each of these op codes is interruptable at the end of each word or byte transfer. If no interrupt requests are active the transfers continue. PC is not incremented to the next op code until the op code is completed. This allows for complete interruptability as long as register integrity is maintained during the interrupt.

BASE OP CODE	MNEMONIC
ØEØØ	MBWU
ØE4Ø	MBWD
ØE8Ø	MBBU
ØECØ	MBBD
ØFØØ	MBWA
ØF4Ø	MBBA
ØF8Ø	MABW
ØFCØ	MABB

\* NOTE: These op codes are all in the third microm.

MBWU	MOVE BLOCK OF WORDS UP
FORMAT: FUNCTION:	MBWU SRC, DST The word string beginning with the word addressed by the source register is moved to successively increasing word addresses as specified by the destination register. The source and destination registers are each incremented by two after each word is transferred. RØ is decremented by one after each transfer, and transfers continue until RØ = $\emptyset$ .
MBWD	MOVE BLOCK OF WORDS DOWN

FORMAT: FUNCTION:

MBWD SRC, DST

The word string beginning with the word addressed by the source register is moved to successively

decreasing word addresses as specified by the destination register. The source and destination registers are each decremented by two after each word is transferred. RØ is decremented by one after each transfer, and transfers continue until  $R\emptyset = \emptyset$ .

INDICATORS:

Unchanged

MOVE BLOCK OF BYTES UP MBBU

FORMAT: FUNCTION: MBBU SRC, DST The byte string beginning with the byte addressed by the source register is moved to successively increasing byte addresses as specified by the destination.

register. The source and destination registers are each incremented by one after each byte is transferred. RØ is decremented by one after each transfer,

and transfers continue until  $R\emptyset = \emptyset$ .

INDICATORS:

Unchanged.

MOVE BLOCK OF BYTES DOWN MBBD

FORMAT:

FUNCTION:

MBBD SRC, DST

The byte string beginning with the byte addressed by the source register is moved to successively decreasing byte addresses as specified by the destination register. The source register, destination register, and RØ, are each decremented by one after each byte is

transferred. Transfers continue until RØ = Ø.

INDICATORS:

Unchanged

MOVE BLOCK OF WORDS TO ADDRESS MBWA

FORMAT:

MBWA

FUNCTION: Same as MBWU except that the destination register is

never incremented.

INDICATORS:

Unchanged

MOVE BLOCK OF BYTES TO ADDRESS **MBBA** 

FORMAT:

SRC, DST MBBA

Same as MBBU except that the destination register is FUNCTION:

never incremented.

INDICATORS:

Unchanged

MOVE ADDRESS TO BLOCK OF WORDS MABW

FORMAT:

MABW SRC, DST

Same as MBWU except that the source register is never FUNCTION:

incremented.

INDICATORS:

Unchanged

MOVE ADDRESS TO BLOCK OF BYTES MABB

SRC, DST

FORMAT:

FUNCTION: Same as MBBU except that the source register is never

incremented.

INDICATORS:

Unchanged

#### FORMAT 9 OP CODES

DOUBLE OPS - ONE OR TWO WORDS - SMØ, DMØ to DM7

<u>15</u> 9	88	6 5	3 2	Ø
OPC	S REG	D M	ODE	D REG

There are 8 op codes in this class representing op codes "7000" to "7FFF". Source mode  $\emptyset$  addressing only is allowed, but destination modes  $\emptyset$  - 7 are allowed for all op codes except 3: JSR and LEA with DMØ will cause an illegal instruction format trap (see chapter 2), and SOB is a special format unique to itself. It is included here only because its destination field is 6 bits long. SOB is a branch instruction. Its 6 bit destination field is a positive word offset from PC, which points to the op code that follows, backwards to the desired address. Forward branching is not allowed. SOB is always a one word op code, and it is used for fast loop control. All other op codes are one word long for DMØ to DM5 addressing and two words long for DM6 or DM7 addressing. The rules for PC relative addressing with DM6 or DM7 are the same as they are for the format 7 op codes. Preliminary decoding of all these op codes except SOB presets the indicator flags as follows: N = 1,  $Z = \emptyset$ ,  $V = \emptyset, C = 1.$ 

BASE OP CODE	MNEMONIC
7øøø	JSR
72ØØ	LEA
7 <b>4</b> ØØ	ASH
76ØØ	SOB
78ØØ	XCH
7AØØ	ASHC
7cøø	MUL
7EØØ	DIV
JSR	JUMP TO SUBROUTINE
FORMAT:	JSR REG, DST
OPERATION:	↓ SP, @SP ← REG
	REG ← PC
	PC ← DST
FUNCTION:	The linkage register is pushed onto the stack; PC,
	which points to the op code that follows, is placed
	in the linkage register; and the destination add-
	ress is placed in PC. DMØ is illegal. The assem-
	bler recognizes the format "CALL DST" as being
•	equivalent to "JSR PC, DST".
INDICATORS:	Preset
•	
LEA	LOAD EFFECTIVE ADDRESS
TODATE	TRANS PRO POR
FORMAT:	LEA REG, DST
OPERATION:	REG ← DST

The destination address is placed into the source FUNCTION:

register. DMØ is illegal. The assembler recognizes

the format "JMP DST" as being equivalent to "LEA PC, DST".

Preset INDICATORS:

EXCHANGE XCH

FORMAT: OPERATION: XCH REG, DST REG **₹**(DST)

FUNCTION:

The source register and destination contents are

exchanged with each other.

INDICATORS:

Preset

SUBTRACT ONE AND BRANCH (IF  $\neq \emptyset$ ) SOB

FORMAT: OPERATION:

REG, DST SOB REG ← REG - 1

FUNCTION:

IF REG  $\neq \emptyset$ , PC  $\leftarrow$  PC - (OFFSET \*2) The source register is decremented by one. If the result is not zero then twice the value of the des-

tination offset is subtracted from PC.

INDICATORS:

Unchanged

ARITHMETIC SHIFT ASH

FORMAT:

ASH REG, DST

FUNCTION:

The source register is shifted arithmetically with the number of bits and direction specified by the destination operand. If (DST) =  $\emptyset$  no shifting occurs. If (DST) = -X then REG is shifted right arithmetically X bits as in an SSRA. If (DST) = +X then REG is shifted left arithmetically X bits as in an SSLA. Only an 8 bit destination operand is used. Thus, DST is a byte For DMØ only the lower byte of the destinaddress.

ation register is used.

INDICATORS:

Preset if (DST) =  $\emptyset$  . Otherwise:

N = Set if REG bit 15 is set

 $Z = Set if REG = \emptyset$ 

V = Set to exclusive or of N and C flags

C = Set to the value of the last bit shifted out of REG

ARITHMETIC SHIFT COMBINED ASHC

FORMAT:

REG, DST ASHC

FUNCTION:

Exactly the same as ASH except that the shift is done

on REG+1: REG. All other comments apply.

Preset if (DST) =  $\emptyset$ . Otherwise: INDICATORS:

N = Set if REG+1 bit 15 is set

 $Z = Set if REG+1: REG = \emptyset$ 

V = Reset

C = Set to the value of the last bit shifted out

MUL MULTIPLY

FORMAT:

MUL REG, DST

OPERATION:

REG+1:REG ← REG \*(DST)

FUNCTION:

An unsigned multiply is performed on the source register and the destination operand. The unsigned

32 bit result is placed in REG+1: REG.

INDICATORS:

N = Set if REG+1 bit 15 is set

 $Z = Set if REG+1: REG = \emptyset$ 

V = Reset

C = Indeterminate

DIV

DIVIDE

FORMAT:

DIV REG, DST

OPERATION:

REG ← [REG+1:REG/(DST)]

REG+1 ← REMAINDER

FUNCTION:

An unsigned divide is performed on the 32 bit source operand REG+1:REG and the destination operand. The unsigned result is placed in REG, and the unsigned remainder is placed in REG+1.No divide occurs and the V-flag is set if REG+1 is greater than or equal to (DST) since the result will not fit into 16 bits. If the

divisor is zero both the V and C flags are set.

INDICATORS:

If no division error:

N = Set if REG bit 15 is set

 $Z = Set if REG = \emptyset$ 

V = Reset

C = Indeterminate
If division error:

N = Reset
Z = Reset
V = Set

 $C = set if (DST) = \emptyset$ 

#### FORMAT 10 OP CODES

DOUBLE OPS - ONE TO THREE WORDS - SMØ TO SM7, DMØ TO DM7.

15 12	11	9	8	6	5	3	2 9	1
OPC	S	MODE	S	REG	D 1	MODE	D REG	ſ

There are 12 op codes in this class representing op codes "1000" to "6FFF" and "9000" to "EFFF". Nine of the op codes are word ops. Three are byte ops. Full source and destination mode addressing with any register is allowed. A one word op code is generated for SM0-SM5 and DM0-DM5 addressing. A two word op code is generated for either SM6-SM7 or DM6-DM7 addressing, but not both. For both SM6-SM7 and DM6-DM7 addressing a three word op code is generated. For a two word op code with word #1 at location X: X + 2 contains the source or destination offset and PC = X + 4 if PC is the register that applies to the offset in location X + 2. For a three word op code with word #1 at location X: X + 2 contains the source offset and X + 4 contains the destination offset. If the source register is PC then PC = X + 4 when added to the offset to compute the source address. If the destination register is PC then PC = X + 6 when added to the offset to compute the destination address.

BASE OP CODE	MNEMONIC
1ØØØ	ADD
2ØØØ	SUB
3øøø	AND
4øøø	BIC
5øøø	BIS
6øøø	XOR
9 <b>øøø</b>	CMP
A <b>ø</b> øø	${\tt BIT}$
вøøø	VOM
CØØØ	CMPB
døøø	MOVB
eøøø	BISB

### WORD OPS

ADD	ADD
FORMAT:	ADD SRC, DST
OPERATION:	$(DST) \leftarrow (SRC) + (DST)$
FUNCTION:	The source and destination operands are added together, and the sum is placed in the destination.
INDICATORS:	N = Set if (DST) bit 15 is set Z = Set if (DST) = $\emptyset$ V = Set if both operands were of the same sign and
	the result was of the opposite sign
	<pre>C = Set if a carry is generated from bit 15 of the result</pre>

SUB SUBTRACT SRC, DST FORMAT: SUB  $(DST) \leftarrow (DST) - (SRC)$ OPERATION: The two's complement of the source operand is added FUNCTION: to the destination operand, and the sum is placed in the destination. INDICATORS: N = Set if (DST) bit 15 is set  $Z = Set if (DST) = \emptyset$ V = Set if operands were of different signs and the sign of the result is the same as the sign of the source operand C = Set if a borrow is generated from bit 15 of the result AND AND FORMAT: AND SRC, DST OPERATION: (DST)  $\leftarrow$  (SRC)  $\triangle$  (DST) FUNCTION: The source and destination operands are logically ANDED together, and the result is placed in the destination. INDICATORS: N = Set if (DST) bit 15 is set  $Z = Set if (DST) = \emptyset$ V = Reset C = Unchanged BIC BIT CLEAR SRC, DST FORMAT: BIC  $(DST) \leftarrow (SRC)_{\Delta}(DST)$ OPERATION: FUNCTION: The one's complement of the source operand is logically ANDED with the destination operand, and the result is placed in the destination. N = Set if (DST) bit 15 is setINDICATORS:  $Z = Set if (DST) = \emptyset$ V = Reset C = Unchanged BIS BIT SET FORMAT: BIS SRC, DST OPERATION: (DST)  $\leftarrow$  (SRC)  $\nabla$  (DST) FUNCTION: The source and destination operands are logically ORED, and the result is placed in the destination. INDICATORS: N = Set if (DST) bit 15 is set $Z = Set if (DST) = \emptyset$ V = Reset C = Unchanged

XOR EXCLUSIVE OR

FORMAT: XOR SRC, DST

OPERATION: (DST)  $\leftarrow$  (SRC)  $\checkmark$  (DST)

FUNCTION: The source and destination operands are logically EX-CLUSIVE ORED, and the result is placed in the destination. INDICATORS:

N = Set if (DST) bit 15 is set

 $Z = Set if (DST) = \emptyset$ 

V = Reset
C = Unchanged

CMP

COMPARE

FORMAT:

CMP SRC, DST

OPERATION:

(SRC) - (DST)

FUNCTION:

The destination operand is subtracted from the

source operand, and the result sets the indicators.

Neither operand is altered.

INDICATORS:

N = Set if result bit 15 is set

 $Z = Set if result = \emptyset$ 

V = Set if operands were of opposite sign and the sign of the result is the same as the sign of (DST) C = Set if a borrow is generated from bit 15 of the

result

BIT

BIT TEST

FORMAT:

BIT SRC, DST

OPERATION:

(SRC)  $\Lambda$  (DST)

FUNCTION:

The source and destination operands are logically ANDED, and the result sets the indicators. Neither

operand is altered.

INDICATORS:

N = Set if result bit 15 is set

 $Z = Set if result = \emptyset$ 

V = Reset
C = Unchanged

VOM

MOVÉ

THOTH

MOV SRC, DST

OPERATION:

(DST)  $\neq$  (SRC)

FUNCTION:

FORMAT:

The destination operand is replaced with the source

operand.

INDICATORS:

N = Set if (DST) bit 15 is set

 $Z = Set if (DST) = \emptyset$ 

V = Reset

C = Unchanged

#### BYTE OPS

For SMØ addressing only the lower byte of the source register is used as an operand. For SMI-SM7 addressing only the addressed memory byte is used as an operand. For DMØ addressing only the lower byte of the destination register is used as an operand with one exception: MOVB will extend the sign through bit 15. For DMI-DM7 addressing only the addressed memory byte is used as an operand.

CMPB

COMPARE BYTE

FORMAT: OPERATION:

CMPB SRC, DST (SRC)<sub>B</sub> - (DST)<sub>B</sub>

FUNCTION: The destination operand is subtracted from the

source operand, and the result sets the indicat-

ors. Neither operand is altered.

INDICATORS: N = Set if result bit 7 is set

 $Z = Set if result = \emptyset$ 

V = Set if operands were of different signs and the sign of the result is the same as the sign

of (DST)B.

C = Set if a borrow is generated from result bit 7

MOVE BYTE

FORMAT: MOVB SRC, DST OPERATION:  $(DST)_B \leftarrow (SRC)_B$ 

FUNCTION: (DSI)B + (SRC)B

FUNCTION: The destination operand is rep

The destination operand is replaced with the source

operand. If DMØ the sign bit (bit 7) is replicat-

ed through bit 15.

INDICATORS: N = Set if (DST)B bit 7 is set

 $Z = Set if (DST)_B = \emptyset$ 

V = Reset
C = Unchanged

BISB BIT SET BYTE

FORMAT: BISB SRC, DST

OPERATION:  $(DST)_{B} \leftarrow (SRC)_{B} \nabla (DST)_{B}$ 

FUNCTION: The source and destination operands are logically

ORED, and the result is placed in the destination.

INDICATORS:  $N = Set if (DST)_B bit 7 is set$ 

 $Z = Set if (DST)_B = \emptyset$ 

V = Reset

C = Unchanged

When using auto increments or decrements in either the source or destination (or both) fields the user must remember the following rule: All increments or decrements in the source are fully completed before any destination decoding begins even if the same index register is used in both the source and destination. The two fields are totally independent.

#### FORMAT 11 OP CODES

DOUBLE OPS - ONE WORD - FLOATING POINT.

_	15 12	11 8	7	6	4	3	2	Ø
Γ	1111	OPC	I	SRC		I	DST	

There are 16 OP Codes in this class representing OP Codes "FØØØ" to "FFFF". Only five are currently defined. They reside in the third microm along with the Format 8 OP Codes. The remaining 11 OP Codes are mapped to the fourth microm for future expansion or customized user OP Codes. All are one word long. Two source and destination addressing modes are available. These two modes, FPØ and FP1, are unique to these OP Codes. Each consists of a 3-bit Register Designation and a 1 bit indirect flag preceding the register designator. For FPØ the indirect bit is Ø, and FP1 it is one. Both the source and destination fields have both addressing modes. The modes are defined as follows:

FPØ The designated register contains the address of the operand.

FPl The designated register contains the address of the address of the operand.

FPØ is the same as standard addressing mode 1, and FP1 is the same as standard addressing mode 7 with an offset of zero.

The computed address is the address of the first word of a 3 word floating point operand. The first word contains the sign, exponent, and high byte of the mantissa. The next higher address contains the middle two bytes of the mantissa, and the next higher address after that contains the lowest two bytes of the mantissa. This format is half way between single and double precision floating point formats, and it represents the most efficient use of microprocessor ROM and register space. The complete format is as follows:

- 1. A 1 bit sign for the entire number which is zero for positive.
- 2. An 8-bit base-two exponent in excess-128 notation with a range of +127, -128. The only legal number with an exponent of -128 is true zero (all zeros).
- A 40 bit mantissa with the MSB implied.

Since every operand is assumed to be normalized upon entry and every result is normalized before storage in the destination addresses, and since a normalized mantissa has a MSB equal to one, then only 39 bits need to be stored. The MSB is implied to be a one, and the bit position it normally occupies is taken over by the exponent to increase its range by a factor of two. The full format of a floating point operand is a follows:

	_	15	14	7	6		Ø
LOCATION	<b>x</b> :	S	EXPONENT		MANTISSA	(HIGH)	
	•	15		8 7			ø
LOCATION	X+2:		MANTISSA		(MIDDLE)		
		15		8 7			Ø
LOCATION	X+4:		MANTISSA		(LOW)		

True zero is represented by a field of 48 zeroes. In effect, the CPU considers any number with an exponent of all zeroes (-128) to be a zero during multiplication and division. For add and subtract the only legal number with an exponent of -128 is true zero. All others cause erroneous results. No registers are modified by any Format 11 OP Code. However, to make room internally for computations 4 registers are saved in memory locations "30" - "38" during the exclution of FADD, FSUB, FMUL and FDIV. These registers are retrieved at the completion of the OP Codes. The registers saved are: the destination address, SP, PC and RØ. No Format 11 OP Code is interruptable (for obvious reasons). FMUL uses location "38" for temporary storage of partial results.

#### FLOATING POINT ERROR TRAPS

Location "3E" is defined as the floating point error trap PC. Whenever an overflow, underflow, or divide by zero occurs a standard trap call is executed with PS and PC pushed onto the stack, and PC fetched from location "3E". I2 is not altered. The remaining memory locations that are reserved for the floating point option ("3A and "3C") are not currently used. The status of the indicator flags and destination addresses during the 3 trap conditions are defined as follows:

### FOR UNDERFLOW (FADD, FSUB, FMUL, FDIV)

N = 1 Destination contain	ins all ze	roes
---------------------------	------------	------

 $z = \emptyset$  (true zero).

V = 1

 $C = \emptyset$ 

#### FOR OVERFLOW (FADD, FSUB, FMUL)

 $N = \emptyset$  Destination not altered in any way.

 $z = \emptyset$ 

V = 1

 $C = \emptyset$ 

#### FOR OVER FLOW (FDIV)

 $N = \emptyset$  Destination not altered if overflow detected

 $Z = \emptyset$  during exponent computation. Undefined

V = 1 otherwise. (Used to save unnormalized

 $C = \emptyset$  partial results during a divide).

### FOR DIVIDE BY ZERO (FDIV)

N = 1 Destination not altered in any way.

 $z = \emptyset$ 

v = 1

C = 1

#### RESERVED TRAPS

If the third microm is in the system and the fourth is not then the last 11 floating point OP codes are the only ones that will cause a reserved OP code trap if executed. If the third microm is not in the system then all Format 8 and 11 OP Codes will cause a reserved OP code trap if executed. However, since the Format 8 OP Codes are interrupt—

able the PC is not advance until the completion of the moves. In all other cases PC is advanced when the OP Code is fetched. For these reasons the PC that is saved onto the stack will point to the offending OP Code during a reserved OP Code trap if and only if the offending OP Code is a Format 8 OP Code. For the Format 11 OP Codes the saved PC will point to the OP Code that follows the offending OP Code. If the user wishes to identify which OP Code caused the reserved OP Code trap he must not preceed a Format 8 OP Code with a Format 11 OP Code or a literal that looks like a Format 11 OP Code.

BASE OP CODE	MNEMONIC
F <b>ØØ</b> Ø	FADD
F1ØØ	FSUB
F2ØØ	FMUL
F3ØØ	FDIV
F4ØØ	FCMP
F5ØØ	
F6ØØ	
F7ØØ	
F8ØØ	
F9 <b>Ø</b> Ø	
FAØØ	
FBØØ	
FCØØ	
FDØØ	
feØØ	
ffØØ	
FADD	FLOATING POINT ADD
FORMAT:	FADD SRC, DST
OPERATION:	$(DST) \leftarrow (DST) + (SRC)$
FUNCTION:	The source and destination operands are added
	together, normalized, and the result is stored
	in place of the destination operand.
INDICATORS:	(if no errors)
	N = Set if the result sign is negative (set).
	Z = Set if the result is zero
	V = Reset
	C = Reset
FSUB	FLOATING POINT SUBTRACT
FORMAT:	FSUB SRC, DST
OPERATION:	$(DST) \leftarrow (DST) - (SRC)$
FUNCTION:	The source operand is subtracted from the
	destination operand. The result is normalized
	and stored in place of the destination operand.
WARNING: THIS OP CO	ODE COMPLEMENTS THE SIGN OF THE SOURCE OPERAND IN

MEMORY AND DOES AN FADD.

INDICATORS: (if no errors)

N = Set if the result sign is negative (set)

Z = Set if the result is zero.

V = Reset

C = Reset

#### FMUL FLOATING POINT MULTIPLY SRC, DST FMUL FORMAT: $(DST) \leftarrow (DST)$ \*(SRC) OPERATION: The source and destination operands are multi-FUNCTION: plied together, normalized, and the result is stored in place of the destination operand. (if no errors) INDICATORS: N = Set if the sign of theresult is negative (set). Z = Set if the result is zero V = Reset C = Reset FLOATING POINT DIVIDE FDIV SRC, DST FDIV FORMAT: $(DST) \leftarrow (DST) / (SRC)$ OPERATION: The destination operand is divided by the source FUNCTION: operand. The result is normalized and stored in place of the destination operand. (if no errors) INDICATORS: N = Set if the sign of the result is negative (set): Z = Set if the result is zero V = Reset C = Reset FLOATING POINT COMPARE FCMP SRC, DST FCMP FORMAT: (SRC) - (DST)OPERATION: The destination operand is compared to the source FUNCTION: operand, and the indicators are set to allow a SIGNED conditional branch.

\*NOTE: True if first words of both operands are not equal.

INDICATORS:

CAUTION: The same physical operand may be used as both the source and destination operand for any of the above floating point OP Codes with no abnormal results except two. They are:

- 1) If an error trap occurs the operand will probably be altered.
- 2) An FSUB gives an answer of -2x, if  $x \neq \emptyset$ , instead of  $\emptyset$ .

N = Set if result is negative

V = Set if arithmetic underflow occurs.\*
C = Set if a borrow is generated. \*

Z = Set if result is zero

# APPENDIX A

# NUMERIC OP CODE TABLE

OP CODE			MNEMONIC	
øøøø	øøøø	øøøø	øøøø	NOP
øøøø	øøøø	øøøø	øøø1	RESET
øøøø	øøøø	øøøø	øø1ø	IEN
øøøø	øøøø	øøøø	ØØ11	IDS
øøøø	øøøø	øøøø	Ø1ØØ	HALT
øøøø	øøøø	øøøø	Ø1Ø1	XCT
øøøø	øøøø	øøøø	<b>Ø</b> 11Ø	BPT
gøøø	øøøø	øøøø	Ø111 Ø111	WFI
g g g g	ØØØØ	ØØØØ	løøø	RSVC
døøø	ØØØØ	øøøø	1ØØ1	
gøøø g	gggg	ØØØØ	1Ø1Ø	RRTT
aaaa Aaaa	gggg ggggg		1Ø11	SAVE
addd Addd		øøøø		SAVS
	ØØØØ	øøøø	11ØØ	REST
gaga gagaa	øøøø	ØØØØ	1101	RRTN
gggg Salaa	øøøø	øøøø	111ø	RSTS
g g g g	øøøø	ØØØØ	1111	RTT
gøøø	øøøø	øøø1	ØREG	IAK
øøø	øøøø	ØØØ1	1REG	RTN
øøø	øøøø	ØØ1Ø	ØREG	MSKO
øøø	øøøø	ØØ1Ø	lreg	PRTN
øøø	øøøø	ØØ11	ARGU	LCC
øøø	øøøø	Ølar	GUME	SVCA
øøø	øøøø	1ØAR	GUME	SVCB
øøø	øøøø	11AR	GUME	SVCC
øøø	øøø1	DISP	LACE	BR
øøø	øø1ø	DISP	LACE	BNE
øøø	ØØ11	DISP	LACE	BEQ
øøø	Ø1ØØ	DISP	LACE	BGE
gøø	ØlØl	DISP	LACE	BLT
øøø	Ø11Ø	DISP	LACE	BGT
gøø	Ø111	DISP	LACE	BLE
gøø	1ØØR	EGØØ	VALU	ADDI
gøø	løør	EGØ1	VALU	SUBI
gøø	1ØØR	EG1Ø	VALU	BICI
gøø	1ØØR	EG11	VALU	MOVI
løøø	1ø1ø	øøмо	DREG	ROR
gøø	1Ø1Ø	ØlMO	DREG	ROL
gøø	1Ø1Ø	1ØMO	DREG	TST
løøø	1Ø1Ø	11MO	DREG	ASL
løøø	1Ø11	ØØMO	DREG	SET
øøø	1Ø11	ØlMO	DREG	CLR
ØØØ	1Ø11	1ØMO	DREG	ASR
løøø	1Ø11	11MO	DREG	SWAB
ggg	11øø	ØØMO	DREG	COM
øøø	11øø	Ø1MO	DREG	NEG
ØØØ	1100	1ØMO	DREG	INC
ØØØ	11ØØ	11MO	DREG	DEC

OP CO	DE			MNEMON	IC
øøøø	11Ø1	ØØMO	DREG	IW2	
øøøø	11Ø1	ØlMO	DREG	SXT	1
øøøø	11Ø1	1ØMO	DREG	TCA	LL
øøøø	11Ø1	11MO	DREG	TJM	5
øøøø	$111\mathscr{D}$	ØØSR	CDST	MBW	IJ
øøøø	111ø	Ølsr	CDST	MBWI	ס
øøøø	111ø	1ØSR	CDST	MBB	IJ
øøøø	111Ø	llsr	CDST	MBB	D
øøøø	1111	øøsr	CDST	MBW.	A
øøøø	1111	Ølsr	CDST	MBB.	A
øøøø	1111	1ØSR	CDST	MAB	N
øøøø	1111	11SR	CDST	MAB:	В
øøø1	SRCR	EGDS	TREG	ADD	
øø1ø	SRCR	EGDS	TREG	SUB	
øø11	SRCR	EGDS	TREG	AND	
ø1øø	SRCR	EGDT	TREG	BIC	
Ø1Ø1	SRCR	EGDT	TREG	BIS	
Ø11Ø	SRCR	EGDS	TREG	XOR	
Ø111	ØØØR	RRDS	TREG	JSR	
Ø111	ØØlR	RRDS	TREG	LEA	
Ø111	ØlØR	RRDS	TREG	ASH	
Ø111	ØllR	RROF	FSET	SOB	
Ø111	1ØØR	RRDS	TREG	XCH	
Ø111	1Ø1R	RRDS	TREG	ASH	3
Ø111	11ØR	RRDS	TREG	MUL	
Ø111	111R	RRDS	TREG	DIV	
1øøø	øøøø	DISP	LACE	BPL	
1ØØØ	øøø1	DISP	LACE	BMI	
1øøø	øø1ø	DISP	LACE	BHI	
1øøø	ØØ11	DISP	LACE	BLO	3
1øøø	ØlØØ	DISP	LACE	BVC	
1øøø	Ø1Ø1	DISP	LACE	BVS	
1øøø	Ø11Ø	DISP	LACE		, BHIS
1ØØØ	Ø111	DISP	LACE		, BLO
1øøø	1ØØR	EGØØ	VALU	SSR	
1øøø	1ØØR	EGØ1	VALU	SSL	
1øøø	1ØØR	EG1Ø	VALU	SSR	
1øøø	1ØØR	EG11	VALU	SSL	
1ØØØ	1Ø1Ø	ØØMO	DREG	ROR	
1øøø	1Ø1Ø	ØlMO	DREG	ROLI	
1øøø	1Ø1Ø	1ØMO	DREG	TST	
1øøø	1Ø1Ø	11MO	DREG	ASL	
1øøø	1011	ØØMO	DREG	SET	
1øøø	1Ø11	ØlMO	DREG	CLR	
1ØØØ	1Ø11	1ØMO	DREG	ASR	
1ØØØ	1Ø11	11MO	DREG	SWA	
1øøø	11ØØ	ØØMO	DREG	COM	
1ØØØ	1100	Ø1MO	DREG	NEG	
1ØØØ	11øø	1ØMO	DREG	INC	
1ØØØ	11ØØ	11MO	DREG	DEC	3

OP CC	DE		<del></del>	MNEMONIC
1 <i>ddd</i>	11ø1	ddwo	DDEC	
1ØØØ	11Ø1	ØØMO	DREG	LSTS
1øøø	•	ØlMO	DREG	SSTS
1ØØØ	11Ø1	1ØMO	DREG	ADC
1øøø	1101	11MO	DREG	SBC
1øøø	111R	EGØØ	VALU	SDRR
1ØØØ	111R	EGØ1	VALU	SDLR
1ØØØ	111R	EG1Ø	VALU	SDRA
1ØØØ	111R	EG11	VALU	SDLA
1ØØ1	SRCR	EGDS	TREG	CMP
1Ø1Ø	SRCR	EGDS	TREG	BIT
1Ø11	SRCR	EGDS	TREG	MQV
11øø	SRCR	EGDS	TREG	CMPB
11Ø1	SRCR	EGDS	TREG	MOVB
111ø	SRCR	EGDS	TREG	BISB
1111	øøøø	ISRC	IDST	FADD
1111	øøø1	ISRC	IDST	FSUB
1111	øø1ø	ISRC	IDST	FMUL
1111	ØØ11	ISRC	IDST	FDIV
1111	Ø1ØØ	ISRC	IDST	FCMP
1111	Ø1Ø1	ISRC	IDST	
1111	ØllØ	ISRC	IDST	
1111	Ø111	ISRC	IDST	
1111	1øøø	ISRC	IDST	
1111	1ØØ1	ISRC	IDST	
1111	1ø1ø	ISRC	IDST	
1111	1Ø11	ISRC	IDST	
1111	11øø	ISRC	IDST	
1111	11Ø1	ISRC	IDST	
1111	111ø	ISRC	IDST	
1111	1111	ISRC	IDST	

#### APPENDIX B

#### ASSEMBLER NOTES

#### FORMAT 1 OP CODES

All are one word op codes except SAVS which is a two word op code. The second word of the SAVS op code is an absolute value.

### FORMAT 2 OP CODES

All are one word with a 3 bit register argument

### FORMAT 3 OP CODE

A one word op code with a 4 bit numeric argument

#### FORMAT 4 OP CODES

All are one word with a 6 bit numeric argument

### FORMAT 5 OP CODES

All are one word with an 8 bit signed PC relative word displacement. The displacement is relative to op code+2. Maximum displacement from the op code is +128, -127 words.

### FORMAT 6 OF CODES

All are one word with a 3 bit register and a 4 bit numeric argument. The stored numeric argument is a positive number from  $\emptyset$  -"F" that equals the actual numeric argument (1-"1 $\emptyset$ ") minus one.

#### FORMAT 7 OP CODES

All are one word op codes for DMØ - DM5 addressing and two word op codes for DM6 - DM7 addressing. For DM6- DM7 addressing the offset is in the second word. If the index register is PC with DM6 - DM7 the offset is relative to op code+4.

#### FORMAT 8 OP CODES

All are one word with a 3 bit source and a 3 bit destination register argument. The count register is implied to be  $R\emptyset$ .

#### FORMAT 9 OP CODES

All have a 3 bit register argument with a 6 bit destination argument that allows DMØ - DM7 addressing. For DMØ - DM5 a one word op code is generated. For DM6 - DM7 a two word op code is generated with the offset in word two. If the index register is PC with DM6-DM7 then the offset is relative to op code+4.

#### FORMAT 10 OP CODES

All have a 6 bit source and a 6 bit destination argument that allow SMØ - SM7 and DMØ - DM7 addressing. For SMØ - SM5 and DMØ - DM5 combined addressing a one word op code is generated. For SM6-SM7 or DM6 - DM7 but not both a two word op code is generated with the offset in word two. If the field with mode 6 or 7 addressing uses PC as the index register then the offset is relative to the op code + 4. For SM6 - SM7 and DM6 - DM7 combined addressing a 3 word op code is generated. Word two contains the source offset, and word 3 contains the destination offset. For SM6 = SM7 with PC the offset is relative to the op code + 4. For DM6 - DM7 with PC the offset is relative to the op code + 6.

Any autoincrements/decrements in the source are fully completed before any destination decoding begins.

### FORMAT 11 OP CODES

All are one word op codes with a 4 bit source and a 4 bit destination argument. Each argument consists of a 3 bit register argument preceded by a 1 bit indirect argument.

#### APPENDIX C

#### PROGRAMMING NOTES

Several of the op codes and addressing modes have personality peculiarities that the user should be aware of. Most of these can be put to good use in particular situations. This appendix attempts to list most of them.

<u>IEN</u>: This instruction allows one more instruction to begin execution before enabling I2.

<u>IDS</u>: This instruction allows one more instruction to begin execution before disabling I2. IDS is therefore interruptable. If such a situation occurs the status of I2 that is included in the pushed PC will equal  $\emptyset$ .

HALT: There is no halt in the microcode. A selection of options is therefore given that allows the user to define HALT for himself.

#### ADDRESSING MODES

In order to clarify the function of the various addressing modes several programming examples are given. In each case assume that the first word of the op code is at location X.

#### SET RØ

Register RØ is set to all ones.

#### CLR @R2

The memory location pointed to by R2 is cleared to zeros. If R2 contained a " $\emptyset 1 \emptyset \emptyset$ " the memory word address " $\emptyset 1 \emptyset \emptyset$ " would be cleared.

#### INC (R3) +

The memory location pointed to by R3 is incremented by one. R3 is then incremented by 2.

#### DEC (PC)+

Location X + 2 is decremented by one, and program control is advanced to location X + 4. This allows for in-line literals in a program, a method that saves a word of memory in most cases.

### SWAB @(R4)+

If R4 contains a " $\emptyset$ 1 $\emptyset$ 0" and location " $\emptyset$ 1 $\emptyset$ 0" contains a " $\emptyset$ 2 $\emptyset$ 0" then the two bytes in location " $\emptyset$ 2 $\emptyset$ 0" are swapped and R4 is incremented to " $\emptyset$ 1 $\emptyset$ 2".

#### COM - (R5)

R5 is decremented by two. The address specified by the altered R5 is one's complemented.

#### NEG - (PC)

A BOZO no-no since location X is the location negated and program control is again transferred to location X after the negation is completed.

### TST @-(R1)

If R = " $\emptyset1\emptyset4$ " and location " $\emptyset1\emptyset2$ " contains a " $1\emptyset\emptyset\emptyset$ " then the following sequence occurs: (1) R1 is decremented by 2 to " $\emptyset1\emptyset2$ ". (2) The contents of location " $\emptyset1\emptyset2$ " (i.e. " $1\emptyset\emptyset\emptyset$ ") becomes the address of the operand to be tested.

## ROR 4 (R4)

The contents of memory location R4 + 4 is rotated right. R4 is not altered. Word two of this op code contains a 4. Program control is advanced to location X + 4 at the completion of the rotate.

### ROL @6(SP)

The contents of memory location SP + 6 contains the address of the operand to be rotated. Word two of this op code contains a 6. Program control is advanced to location X + 4 at the completion of the rotate.

#### JSR PC, TAG

Location X + 2 contains the byte offset from location "TAG" to location X + 4. The address of location X + 4 is pushed onto the stack, and the address of location "TAG" is placed in PC.

#### JSR R5, TAG

Location X + 2 contains the byte offset from location "TAG" to location X + 4. The content of register R5 is pushed onto the stack, the address of location X + 4 is placed in R5, and the address of location "TAG" is placed in PC.

### JSR PC, (R4)+

Location X + 2 is pushed onto the stack, R4 is moved to PC, and R4 is incremented by two.

#### JSR PC,@(SP)+

This is a co-routine call. Pay attention:

 The contents of the location pointed to by SP is saved in CPU register "TMPA".

- 2) SP is incremented by two.
- 3) The address of location X + 2 is pushed onto the stack
- 4) CPU register "TMPA" is moved to PC

The effect of all this is to swap the top word on the stack with the address of location X + 2 without altering SP or stack size. Consider the following routine.

SUBR: JSR PC,2(PC)
TAGA: JSR PC,<sup>@</sup>(PC)

TAGB:

RTN PC

The first JSR places the address of TAGA on the stack and executes the routine starting at TAGB. The RTN PC transfers control to location TAGA when it is executed. The second JSR places address TAGB onto the stack and into PC, effectively leaving PC unaltered. The second time the RTN PC is executed program control passes to location TAGB. The third time the RTN PC is executed program control passes back to the routine that call subroutine SUBR. Since TAGA and TAGB are never addressed explicitly both of the labels could be eliminated from the program. If left in then the "2(PC)" could be replaced with "TAGB".

#### CMP $(R\emptyset)+$ , $(R\emptyset)+$

If  $R\emptyset = "\emptyset 1 \emptyset \emptyset"$  then the contents of location " $\emptyset 1 \emptyset \emptyset$ " is compared to the contents of location " $\emptyset 1 \emptyset 2$ ", and  $R\emptyset$  is incremented to " $\emptyset 1 \emptyset 4$ ". All source auto increments or decrements are completed before destination decoding begins.

#### MOV @R2,-(R2)

If R2 = " $\emptyset$ 1 $\emptyset$ 6" then the contents of location " $\emptyset$ 1 $\emptyset$ 6" is moved to location " $\emptyset$ 1 $\emptyset$ 4", and R2 is decremented to " $\emptyset$ 1 $\emptyset$ 4".

#### BIT #2,@#4

The contents of absolute memory location 4 is tested against the literal value 2. This is a three word op code with word two containing a 2 and word three containing a 4. This op code works on location 4 from anywhere in memory.

#### CMP (PC)+,TAG

This won't work. The assembler generates a two word op code for this with the destination offset in word two. The execution of the op code, however, uses word two as a literal and word three (which does not exist) as the destination offset. By swapping the source and destinations around then an in-line literal could be used for word three, and word two would contain a valid source offset.

## JSR PC, (PC)+

The address of location X + 4 is pushed onto the stack, and PC gets the address of location X + 2.

#### JSR R5, (PC)+

The contents of R5 are pushed onto the stack, R5 gets the address of location X + 4, and PC gets the address of location X + 2.

### MOVB $(R\emptyset)+$ , $(R\emptyset)+$

If  $R\emptyset = "\emptyset1\emptyset2"$  then the contents of memory byte location " $\emptyset1\emptyset2"$  is moved to memory byte location " $\emptyset1\emptyset3"$ , and  $R\emptyset$  is incremented to " $\emptyset1\emptyset4"$ .

#### MOVB (SP) + R1

The contents of the memory byte addressed by SP is moved to the lower byte of Rl, the sign bit (bit 7) is replicated through bit 15 of Rl, and SP is incremented by 2. SP is always autoincremented or autodecremented by two.

#### CLRB (PC)+

The contents of the lower byte memory location X + 2 is cleared to zeros. The upper byte (X + 3) is not affected. PC is incremented by two. PC is always autoincremented or autodecremented by two.

#### BISB RØ, Rl

The lower bytes of register RØ is logically ORED with the lower byte of register Rl. The upper byte of Rl is not altered.

### MOVB @(R2)+,@-(R3)

If R2 contains a " $\emptyset1\emptyset\emptyset$ " and R3 contains a " $\emptyset2\emptyset\emptyset$ " then location " $\emptyset1\emptyset\emptyset$ " contains the byte address of the source operand and location " $\emptyset1FE$ " contains the address of the destination byte that is to receive the source byte. R2 is incremented by two, and R3 is decremented by two since they point to addresses of (16 bit) addresses.

#### JSR SP, TAG

Not recommended since the value of the stack is lost. Perfectly legal however.

### SAVS and RSTS

Although designed to be used for automatic register and I/O priority level saving and restoring, the lack of hardware priority masking does not alter the operation or the op codes. The SAVS op code is usually the first instruction executed in a device interrupt routine, and the RSTS is the last. The priority mask can use a one bit as an enable or disable with bit  $\emptyset$  the highest or lowest priority level. Such decisions are made by the hardware.

#### POWER FAIL

Two levels of power fail are provided for in the firmware. The hardware may use two, one, or no levels of power fail. The three modes are discussed in increasing order of complexity.

NO LEVELS: External address register bit 7 is hardwired to Ø, and a prayer is offered.

ONE LEVEL: The detection of a power fail sets bit 7 of the external status register and the CPU RESET line. When the power fail disappears the CPU RESET line is reset, but bit 7 of the external status register remains set. The Line Clock Clear State Code (see appendix D) clears bit 7 of the external status register (and bits 5, 6 if used). A system power up is then executed.

TWO LEVELS: This requires two hardware functions, AC LOW and DC LOW, plus two levels of power fail; AC and DC. It all works like this: If AC power begins to deteriorate AC LOW is set first. This sets bit 7 of the external status register and generates an interrupt via  ${\tt I} \emptyset$  or Il. If AC power does not deteriorate too far then nothing else happens except that bit 7 of the external status register is reset when power is restored. If AC power continues to deteriorate then eventually DC power will begin to deteriorate. When this happens DC LOW is set and DC LOW sets CPU RESET. AC LOW is still set and it maintains bit 7 of the external status register. When power is restored DC LOW is reset. resets CPU RESET. A power up sequence is initiated, and the Line Clock Clear State (see appendix D) clears The External Status Register bit 7 (plus 5 and 6 if they are used). If the user wishes to be able to execute a programmed power fail routine even during a sudden and complete power failure then the DC power supply must be strong enough to run the CPU and MEMORY for at least 2 milliseconds. The power fail interrupt must also be programmed, and the interrupts enabled.

The use of the Line Clock Clear State Code to clear bits 5-7 on a CPU RESET function (plus the line clock of course) should have no effect on normal system operation. Should an error occur during a non-vectored interrupt the error would be cleared momentarily and then set again as CPU RESET obviously could not have been generated. If it had been then the system could not be in the non-vectored interrupt routine.

#### PARITY AND BUS ERRORS

These functions are also part of the CPU RESET function along with power fail/up. In order to get only one or the other then bit 7 of the external status register must be reset when the CPU RESET function

is activated. In order to generate a valid CPU RESET the CPU RESET line must be held active for three clock cycles. Longer is fine, but the CPU goes into a wait state until the CPU RESET is reset. If more than one error exists at one time then the highest priority error is the one honored The priority, from highest to lowest, is:

Power Fail Bus Error Parity Error

If all 3 functions are reset a power up is assumed. All 3 functions have a bit associated with them in the external status register. Once set these bits stay set until cleared by the Line Clock Clear State Code (see appendix D) that is generated during the first phases of the reset routine. See chapter two "Power Up Options".

#### APPENDIX D

#### MICROM STATE CODE FUNCTIONS

Below is a list of MICROM STATE CODE FUNCTIONS for the WD1600 with a brief description of what each does. More elaborate descriptions, where necessary, follow the table.

CODE	MNEMONIC	FUNCTION
ØØØ1 ØØ1Ø ØØ11 Ø1ØØ	PMSK RUN IORST INTEN	Priority mask out Macro instruction fetch I/O reset I2 set
Ø1Ø1 Ø11Ø Ø111 1ØØØ	INTDS ESRR SRS BYTE	I2 reset External status register request System reset Read byte operation
1001 1010 1011 1100 1101	RMWW RMWB RLCI EARR	Read-modify-write word Read-modify-write byte Reset line clock interrupt External address register request Duplicate of "BYTE"
111Ø 1111		Duplicate of "RMWW"  Duplicate of "RMWB"

PMSK: The state code is generated on an OUTPUT WORD instruction when a new mask is written into location "2E". It signals the I/O devices that a new interrupt mask is on the DAL.

RUN: Generated during macro instruction fetch for a run light.

<u>IORST</u>: Generated during a RESET macro op code to reset I/O devices to some preset state.

INTEN: Enables the interrupt enable line -I2.

INTDS: Disables the interrupt enable line -I2.

ESRR: Generated during an INPUT STATUS BYTE micro op code to indicate that the external status register is being requested. See note 1.

SRS: Generated during a power up for a master system reset. This code is followed by a 300 cycle wait to allow time for any reset functions the hardware generates to be completed before any DAL requests are generated.

BYTE: Generated during an INPUT BYTE micro op code to indicate a read byte operation without a read-modify-write.

RMWW: Generated during an INPUT WORD micro op code with RMW active to indicate a read-modify-write word sequence.

RMWB: Generated during an INPUT BYTE micro op code with RMW active to indicate a read-modify-write byte sequence.

- RLCI: Generated during a CPU RESET or a non-vectored interrupt without a power fail to clear both the line clock interrupt and external status register bits 5-7.
- EARR: Generated during an INPUT STATUS BYTE micro op code to indicate a request for the external address register during the user bootstrap routine.
- CODES "D" "F": Duplicates of codes "8" "A" respectively except that these codes appear as a part of the READ micro op codes instead of as a part of the INPUT micro op codes. Either or both may be used by the hardware as is convenient. These codes preceed the others. They are generated only once, however, instead of repeating in the event of a wait state as the others do.
- NOTE 1: INPUT STATUS BYTE is not a function of reply and does not generate a SYNC. For these reasons the DAL must be tri-stated if a DMA device also exists. The data is always gated onto the lower byte. The upper byte is ignored.
- NOTE 2: Lack of state codes "8" "A" or "D" "F" during a READ INPUT sequence implies a read word operation without read-modify-write.

where the constant of the constant (x,y) is the constant of the constant (x,y)

#### APPENDIX E

### OP CODE TIMINGS

All times are in cycles. Timings include all OP Code fetches, memory reads, and memory writes applicable to each. Timings assume that the memory is running with full speed with respect to the CPU. This requires a 16 Bit access time = 1 CPU cycle, and a 16 Bit memory read/write cycle time = 2 CPU cycles. One CPU cycle = 300 NS @ 3.3 MHZ, UØØ NS @ 2.5 MHZ, and 500 NS @ 2 MHZ clock rates. Timings are included for SMØ and DMØ as basic with additions as necessary in tables that follow the OP Codes for SM1-7 and DM1-7 timings.

### FORMAT ONE OP CODES

OP CODE	# CYCLES	. •
NOP	1Ø	
RESET	1Ø	
IEN	1Ø	
IDS	1Ø	
HALT	16+	
XCT	44 + OP CODE	EXECUTED
BPT	24	1
WFI	16+	
RSVC	62	*
RRTT	60	-
SAVE	46	
SAVS	65	
REST	48	•
RRTN	52	
RSTS	64	
RTT	13	. *

## FORMAT TWO-FOUR OP CODES

OP CODE	# CYCLES		
IAK	1Ø		
R'I'N	12		
MSKO	1Ø		
PRTN	22		
LCC	4 A 7		
SVCA	37		
SVCB	73		
SVCC	71		

#### FORMAT FIVE OP CODES

All branches = 9 cycles if branch occurs or not.

## FORMAT SIX OP CODES

OP CODE	# CYCLES
ADDI	9
SUBI	9
BICI	9
MOVI	9
SSRR	8 + (5 x # bits shifted)
SSLR	8 + (5 x # bits shifted)
SSRA	8 + (7 x # bits shifted)
SSLA	8 + (5 x # bits shifted)
SDRR	20 + (7 x # bits shifted)
SDLR	20 + (7 x # bits shifted)
SDRA	20 + (9 x # bits shifted)
SDLA	$20 + (7 \times 3 \text{ bits shifted})$

### FORMAT 7 OP CODES - DMØ

OP CODES	# CYCLES		OP CODES	# CYCLES	
ROR	1Ø		RORB	9	
ROL	1Ø		ROLB	9	
TST	1Ø		TSTB	9	
ASL	1Ø		ASLB	9	
SET	1Ø		SETB	1Ø	
CLR	1Ø		CLRB	9	
ASR	12		ASRB	11	
SWAB	1Ø		SWAD	21	
COM	1Ø		COMB	9	
NEG	1Ø		NEGB	9	
INC	1Ø		INCB	9	
DEC	1Ø		DECB	9	
IW2	1Ø		LSTS	15	
SXT	12		SSTS	1Ø	
TCALL	21		ADC	11	
TJMP ·	16		SBC	11	
FOR WORD OPS	AND:		FOR BYTE OPS AN	ND:	
DMl	ADD	4	DMl	ADD	3
DM2	ADD	4	DM2	ADD	3 *
DM3	ADD	8	D <b>M</b> 3	ADD	7
DM4	ADD	6	DM4	ADD	5 *
DM5	ADD	1Ø	DM5	ADD	9
DM6	<b>AD</b> D	1Ø	DM6	ADD	9
DM7	ADD	14	DM7	ADD	13

For DMl - DM7 and:

CLR subtract 1 cycle SWAB subtract 1 cycle

\*NOTE: Add 2 more if SP or PC.

# FORMAT 8 OP CODES

	To provide the section of the sectio			
OP CODE	# CYCLES (ASSUMES NO INTERRUPTS)			
MBWU	17 + (16 X # words moved)			
MBWD	15 + (16 X # words moved)			
MBBU	17 + (15 X # bytes moved)			
MBBD	15 + (15 x # bytes moved)			
MBWA	19 + (16 X # words moved)			
MBBA	19 + (15 X # bytes moved)			
MABW	19 + (16 X # words moved)			
MABB	19 + (15 X # bytes moved)			
	FORMAT 9 OP CODES - DAMO	•		
OP CODE	# CYCLES			
OI CODE	II Ca Cada C			
JSR*	22			
LEA*	15			
ASH	19 if DST = $\emptyset$ ; 22 + (5 X count) if DST> $\emptyset$ ; 25+(7	X count.)	if DST <	ø.
SOB	1Ø if no branch, 13 if branch			
XCH	23			
ASHC	19 if DST = $\emptyset$ ; 38 + (7 X count) if DST> $\emptyset$ ; 38+(9	X count)	if DST <	Ø
MUL	183			
DIA	29 if divisor error, 202 if no divisor error			
*NOTE: DMØ i	illegal. Used as base figure only.			
FOR ALL OP CO	ODES EXCEPT SOB AND:			
DMl add Ø				
DM2 add 2				
DM3 add 2				
DM4 add 2				
DM5 add 4 DM6 add 4				
DM6 add 4 DM7 add 8				
DM7 add 6				

# FORMAT 10 OP CODES - SMO AMD DMO

OP CODE	# CYCLES
ADD	11
SUB	11
AND	11
BIC	11
BIS	11
XOR	11
CMP	. 11
BIT	11
VOM	11
CMPB	11
MOVB	12
BISB	11

For SM1: add 3 for word ops, 1 for byte ops.
For SM2: add 4 for word ops, 2 for byte ops. \*
For SM3; add 7 for word ops, 5 for byte ops.
For SM4; add 5 for word ops, 3 for byte ops. \*
For SM5; add 9 for word ops, 7 for byte ops.
For SM6; add 9 for word ops, 7 for byte ops.
For SM7; add 13 for word ops, 11 for byte ops.

For DM1; add 4 for word ops, 3 for byte ops.
For DM2; add 4 for word ops, 3 for byte ops. \*
For DM3; add 8 for word ops, 7 for byte ops.
For DM4; add 6 for word ops, 5 for byte ops. \*
For DM5; add 10 for word ops, 9 for byte ops.
For DM6; add 10 for word ops, 9 for byte ops.
For DM7; add 14 for word ops, 13 for byte ops.

For MOVB and DM1-DM7 subtract 1 cycle.

\*NOTE: Add 2 if SP or PC

### FORMAT 11 OP CODES - ALL ADDRESSING MODES

FADD:	If exponent difference > 39 Worst Case	:	138-145 638
	Typical	:	180-420
FSUB:	If exponent difference > 39	:	141-148
	Worst Case	:	641
	Typical	:	190-430
FMUL:	If either operand = 0	:	108-111
	Worst Case	:	805
	Typical	:	590-780
FDIV:	If divide by Ø	:	96
	If divide into Ø	:	118
	Worst Case	:	1596
	Typical	:	280-1210
FCMP:		:	49-86